



**MPLS Inter-Carrier Interconnect  
(MPLS-ICI)**

**Technical Specification**

**IP/MPLS Forum 19.0.0**

**IP/MPLS Forum Technical Committee  
April 2008**

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## Table of Contents

<b>1</b>	<b>INTRODUCTION</b> .....	<b>1</b>
<b>2</b>	<b>SCOPE</b> .....	<b>1</b>
<b>3</b>	<b>DEFINITIONS</b> .....	<b>3</b>
<b>4</b>	<b>ACRONYMS</b> .....	<b>3</b>
<b>5</b>	<b>REFERENCES</b> .....	<b>4</b>
5.1	NORMATIVE REFERENCES .....	4
5.2	INFORMATIVE REFERENCES .....	5
<b>6</b>	<b>REFERENCE ARCHITECTURE</b> .....	<b>8</b>
<b>7</b>	<b>PHYSICAL LAYER (LAYER 1)</b> .....	<b>10</b>
<b>8</b>	<b>DATA LINK LAYER (LAYER 2)</b> .....	<b>10</b>
<b>9</b>	<b>MECHANISMS FOR LSP ESTABLISHMENT</b> .....	<b>10</b>
9.1	ALL-STATIC CONFIGURATION .....	10
9.1.1	Resiliency .....	11
9.1.2	MIBs .....	15
9.2	STATICALLY CONFIGURED AND SIGNED LSPs: PWs AND TE-TUNNELS .....	17
9.3	DYNAMICALLY ESTABLISHED LSPs: BGP/MPLS IP VPN MULTI-AS BACKBONES OPTION B, BGP LABELED IPV4 PATHS, AND TE-TUNNELS .....	17
9.4	MIBS .....	17
<b>10</b>	<b>CONNECTION ADMISSION CONTROL</b> .....	<b>17</b>
<b>11</b>	<b>FORWARDING</b> .....	<b>18</b>
11.1	TRAFFIC MANAGEMENT .....	18
11.1.1	Classes of Service and EXP Mappings in the Forwarding Plane .....	18
11.1.2	Traffic Policing .....	19
11.1.3	Traffic Shaping .....	19
11.2	PATH MTU HANDLING AND FRAGMENTATION .....	20
11.2.1	Label stack depth .....	20
11.2.2	Fitting packet sizes to PMTU .....	21
11.2.3	MTU Requirements for MPLS-ICI .....	21
11.3	LOAD BALANCING .....	22
11.4	TIME TO LIVE (TTL) PROCESSING .....	22
11.5	MIBS .....	23
<b>12</b>	<b>OAM</b> .....	<b>24</b>
12.1	CONNECTION VERIFICATION: "ALWAYS ON" DEFECT DETECTION AND HANDLING .....	24
12.2	ON DEMAND DIAGNOSTICS .....	25
12.3	MIBS .....	27
<b>13</b>	<b>SECURITY AND CONFIDENTIALITY</b> .....	<b>28</b>
13.1	CONTROL PLANE PROTECTION .....	28
13.1.1	Authentication of Signaling Sessions .....	28
13.1.2	Protection against DoS attacks in the Control Plane .....	29

13.1.3	Protection against Malformed Packets.....	29
13.1.4	Ability to Enable/Disable Specific Protocols.....	29
13.1.5	Protection against Incorrect Cross Connection.....	29
13.1.6	Protection Against Spoofed Updates and Route Advertisements.....	29
13.1.7	Protection of Confidential Information.....	30
13.2	DATA PLANE PROTECTION.....	30
13.2.1	Protection against DoS in the Data Plane.....	30
13.2.2	Protection against Label Spoofing.....	30
<b>ANNEX A</b>	<b>IP VPN.....</b>	<b>31</b>
A.1	ROUTING.....	31
A.1.1	ROUTE TARGET ALLOCATIONS – MULTI-AS BACKBONES OPTION B.....	31
A.1.2	Route Distinguisher (RD) Allocations – Multi-AS Backbones Option B.....	32
A.1.3	Route Target Advertisement – Multi-AS Backbones Option B.....	33
A.1.4	VPN-IPv4 Route Summarization.....	33
A.2	LABEL VALUE ACCEPTANCE – MULTI-AS BACKBONES OPTION B.....	33
A.3	DSCP/PRECEDENCE-EXP MAPPING.....	33
A.4	RESILIENCY.....	35
A.5	ADMISSION CONTROL POLICY: VPN-IPv4 ROUTE LEARNING RESTRICTION REQUIREMENT.....	36
A.6	MIBS.....	36
A.7	BGP/MPLS IPv4 VPN – MULTI-AS BACKBONES OPTION A (NON-NORMATIVE).....	37
A.8	BGP/MPLS IPv4 VPN – MULTI-AS BACKBONES OPTION B (NON-NORMATIVE).....	38
A.9	INTERCONNECT PERFORMANCE MEASUREMENT (NON-NORMATIVE).....	39
<b>ANNEX B</b>	<b>LABELED IPV4 ROUTES.....</b>	<b>40</b>
B.1	MP-BGP – LABEL DISTRIBUTION.....	40
B.2	ROUTING.....	40
B.3	RESILIENCY.....	40
B.4	POLICIES AND ADMISSION CONTROL.....	41
B.5	MIBS.....	41
<b>ANNEX C</b>	<b>PSEUDOWIRES.....</b>	<b>42</b>
C.1	STATICALLY-CONFIGURED AND DYNAMICALLY SIGNED MPLS PSEUDOWIRES.....	42
C.1.1	Signaling.....	42
C.1.2	Connection Admission Control.....	43
C.1.3	Routing.....	43
C.1.4	Resiliency.....	43
C.2	PW OAM.....	48
C.2.1	PW OAM Mechanisms.....	49
C.2.1.1	PW Connection Verification: “always on” failure detection and notification.....	49
C.2.1.2	Diagnostics.....	51
C.2.2	Failure Detection and Notification Procedures.....	52
C.3	MIBS.....	53
<b>ANNEX D</b>	<b>TRAFFIC ENGINEERING (TE)-TUNNELS.....</b>	<b>55</b>
D.1	STATICALLY-CONFIGURED AND SIGNED TE-TUNNELS.....	55
D.1.1	Signaling.....	55
D.1.2	MIBS.....	55
D.2	DYNAMICALLY ESTABLISHED TE-TUNNELS.....	56
D.2.1	Signaling.....	56
D.2.2	Routing.....	56
D.2.3	Resiliency.....	56
D.2.4	MIBS.....	57
D.3	ADMISSION CONTROL.....	57
D.4	PROTECTION OF CONFIDENTIAL INFORMATION.....	59

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D.5	CLASS TYPE MAPPINGS .....	60
<b>ANNEX E</b>	<b>VOICE OVER IP .....</b>	<b>61</b>
<b>APPENDIX I</b>	<b>INFORMATIVE DESCRIPTION OF MIBS.....</b>	<b>62</b>
I.1	MPLS LSR MIB [RFC3813].....	62
I.2	DIFFERENTIATED SERVICES MIB [RFC3289].....	63
I.3	BI-DIRECTIONAL FORWARDING DETECTION (BFD) MIB [BFD MIB] .....	64
I.4	BGP MPLS L3VPN MIB [RFC4382] .....	65
I.5	MPLS LDP MIB [RFC3815].....	66
I.6	PSEUDOWIRE MIB [PWMIB] .....	67
I.7	PSEUDOWIRE TO MPLS PSN MIB [PWMPLSMIB].....	68
I.8	MPLS TRAFFIC ENGINEERING (TE) MIB [RFC3812].....	69
<b>APPENDIX II</b>	<b>MPLS-ICI FORWARDING BEHAVIOR AND EXP BIT MAPPING CONFIGURATION EXAMPLE 70</b>	
II.1	OVERVIEW .....	70
II.2	EXAMPLE OF BI-LATERAL AGREEMENT FORWARDING BEHAVIOR AND EXP-BIT MAPPINGS .....	71
II.3	ANOTHER EXAMPLE OF BI-LATERAL FORWARDING BEHAVIOR MAPPING.....	73
II.4	EXAMPLE OF CANONICAL FORWARDING BEHAVIOR MAPPING .....	76

## Table of Figures

Figure 1: Control Protocols involved in extending MPLS services across an MPLS-ICI interconnecting two provider networks .....	8
Figure 2: Services and LSPs extended across an MPLS-ICI interconnecting two provider networks.....	9
Figure 3: Protocol Stack .....	9
Figure 4: One-hop MPLS-ICI Protection .....	12
Figure 5: Multi-Hop LSP protection over an MPLS-ICI .....	13
Figure 6: Multi-Hop bypass tunnel (Many:One) protection for the MPLS-ICI.....	14
Figure 7: Inter-provider IP VPN service using RFC 4364 Multi-AS Backbones Option A .....	38
Figure 8: Inter-provider IP VPN service using RFC 4364 Multi-AS Backbones Option B.....	39
Figure 9: Multi-Segment PW Reference Model with a PSN Tunnel at the MPLS-ICI.....	42
Figure 10: Multi-Segment PW Reference Model without a PSN Tunnel at the MPLS-ICI.....	43
Figure 11: ASBR Fault Detection/Notification at the MPLS-ICI used in end-to-end protection (non-normative).....	45
Figure 12: MS-PW Segment Protection over Parallel MPLS-ICIs.....	46
Figure 13: Reference model for a MS-PW across an MPLS-ICI.....	52

## Table of Tables

Table 1 - LSP ping FEC and sub-TLVs.....	26
Table 2 - LSP ping reply modes .....	26
Table 3: Example Usage of EXP $\leftrightarrow$ PHB Mappings at SP-1's ASBR .....	71
Table 4: Example Usage of EXP $\leftrightarrow$ PHB Mappings at SP-2's ASBR .....	72
Table 5: Example Usage of EXP $\leftrightarrow$ PHB Mappings at SP-A's ASBR.....	74
Table 6: Example Usage of EXP $\leftrightarrow$ PHB Mappings at SP-B's ASBR.....	75





## 1 Introduction

This document presents the MPLS Inter-Carrier Interconnect (MPLS-ICI) Technical Specification. The objective of this Technical Specification is to enable a set of MPLS services, listed in Section 2, across administrative domains operated by different carriers. This specification could also be useful to individual carriers operating multiple networks resulting, for example, from mergers and acquisitions.

The MPLS-ICI is a bi-directional IP-MPLS logical link between an Autonomous System Border Router (ASBR) in one service provider (SP) network and an ASBR in another service provider network. The logical link appears internally to each ASBR as a single logical interface. There may be more than one MPLS-ICI between a pair of ASBRs. An ASBR, as specified in this document, is not required to support IP forwarding of user datagrams when the service does not require it.

One or more Label Switched Paths (LSPs) may cross the MPLS-ICI. A given ASBR will terminate an LSP that crosses the MPLS-ICI, or it will switch the LSP toward the destination.

This Technical Specification focuses on mechanisms that enable the establishment of inter-carrier LSPs, including pairs of LSPs that form pseudowires (PW). It covers configuration, routing, signaling, and forwarding as well as management of the label switched paths.

This document is organized into a base document and services annexes. The base document describes the base functional elements of the MPLS-ICI. It also describes the protocols, information elements and procedures required to provide the base functional elements. This information applies to all MPLS services that utilize the MPLS-ICI. The scope section lists the MPLS services covered in this document.

Each optional MPLS-service annex describes the protocol information content, and procedures for an MPLS service that uses the MPLS-ICI. An MPLS-service annex describes the service, the use of the base MPLS-ICI elements, any additional attributes conveyed in the MPLS-ICI messages, if needed, and procedural processing of the attributes specific to the service.

This document also specifies the Management Information Base (MIB) modules as standardized by the IETF to configure and manage the various aspects of an MPLS-ICI service in the appropriate sections. Compliance groups from the SNMP MIBs are referenced using the descriptive name and specific MIB in this document. An ASBR should support the management information model consistent with these groups. An informative appendix for each major MIB provides a summary description of what each MIB group or notification performs.

Although the document uses MIBs to describe the management information model, an ASBR may implement various management interfaces for read-only or read/write support such as CLI, XML, enterprise MIBs, or configuration files besides SNMP.

## 2 Scope

The purpose of this document is to specify capabilities that enable inter-carrier MPLS services. This document is targeted to equipment vendors to specify the necessary ASBR requirements and to service providers to provide guidance on how to use these capabilities. This document uses protocols developed in appropriate standard bodies (e.g., IETF, ITU) where applicable.

Following are the MPLS-ICI services covered in this document:

- 
- Border Gateway Protocol (BGP)/MPLS IP Virtual Private Networks (VPNs) Multi-AS Backbones Option A [RFC4364]
  - BGP/MPLS IP VPNs Multi-AS Backbones Option B [RFC4364]
  - Inter-domain tunnel establishment using BGP for IP version 4 (IPv4) [RFC3107]
  - PWs and multi-segment pseudowires (MS-PW) (statically signaled and routed, statically routed and stitched with signaled segments) [RFC3995] [MS-PW\_Requirements]
  - Inter-domain Traffic Engineered (TE) tunnels (statically signaled and routed, dynamically signaled and routed) [RFC5151]

For the services addressed in this document, following are within the scope of this Technical Specification:

- MPLS forwarding
- IPv4 processing when processing of the IP header is required as a matter of forwarding MPLS-labeled packets
- IPv4 processing for signaling and routing control
- MPLS Pseudowires with static configuration
- Option 'A' of the 'Multi-AS Backbones' section of [RFC4364]. This is a case where the IP traffic forwarded on the MPLS-ICI is unlabeled IP but the MPLS service is a BGP/MPLS IP VPN service.
- TE Tunnels and Resource Reservation Protocol with Traffic Engineering (RSVP-TE) signaling
- Label Distribution Protocol (LDP) signaling for pseudowires
- MPLS IP VPNs
- BGP for inter-domain routing and establishing label switched paths
- MPLS Operation, Administration and Maintenance (OAM)
- Connection Admission Control (CAC)
- Resiliency
- Classes of Service (CoS)
- Policies
- Security

- MIBs: where no previous standards work exists, this document describes a management information model in the form of an Enterprise MIB that a router would need to implement in order to support the MPLS-ICI base function or service.

Following are out of the scope of this Technical Specification:

- Inter-Carrier MPLS-based services that do not involve MPLS label processing at the interconnect for forwarding purposes. Therefore, usage of the MPLS-ICI by non-MPLS services (e.g., IP) is outside the scope of the MPLS-ICI. The only exception is RFC4364 Multi-AS Backbones Option A.
- IP version 6 (IPv6)

### 3 Definitions

**must, shall or mandatory** — the item is an absolute requirement of this Technical Specification.

**should** — the item is desirable.

**may or optional** — the item is not compulsory, and may be followed or ignored according to the needs of the implementer.

### 4 Acronyms

ATM	Asynchronous Transfer Mode
AS	Autonomous System
ASBR	Autonomous System Border Router
BGP	Border Gateway Protocol
MP-BGP	Multi-Protocol BGP
CE	Customer Edge
CoS	Class of Service
DoS	Denial of Service
ERO	Explicit Route Object
IETF	Internet Engineering Task Force
IP	Internet Protocol
ITU-T	International Telecommunication Union - Telecom
LDP	Label Distribution Protocol
LSP	Label Switched Path
LSR	Label Switching Router
MIB	Management Information Base
MPLS	MultiProtocol Label Switching
MS-PW	Multi-segment Pseudowire
PE	Provider Edge

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PHB	Per-Hop Behavior
PSN	Packet Switched Network
PW	Pseudowire
QoS	Quality of Service
RSVP-TE	Resource Reservation Protocol with Traffic Engineering Extensions
SP	Service Provider
S-PE	Switching PE
TE	Traffic Engineering
T-PE	Terminating PE
VPN	Virtual Private Network

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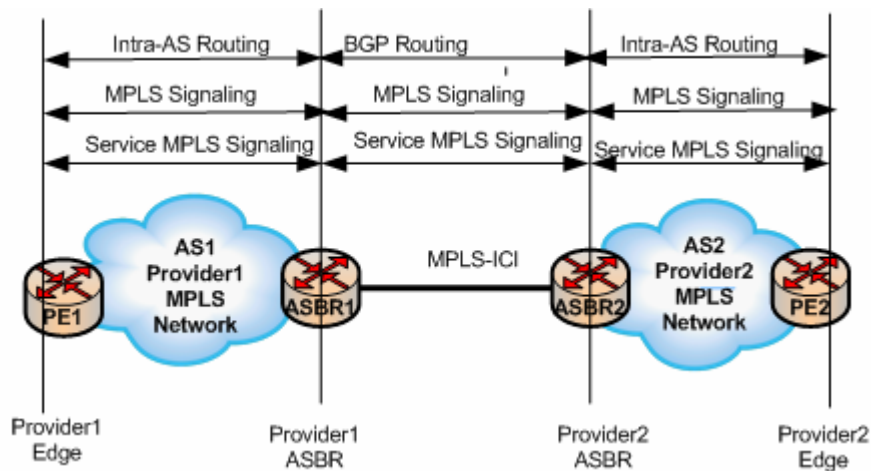
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## 6 Reference Architecture

Following are reference model diagrams that describe the architecture of the MPLS inter-carrier interconnection (MPLS-ICI), illustrating the control protocols and service extensions across an inter-carrier MPLS interface, and the protocol stack. These reference models are applicable to all MPLS services discussed in this document.

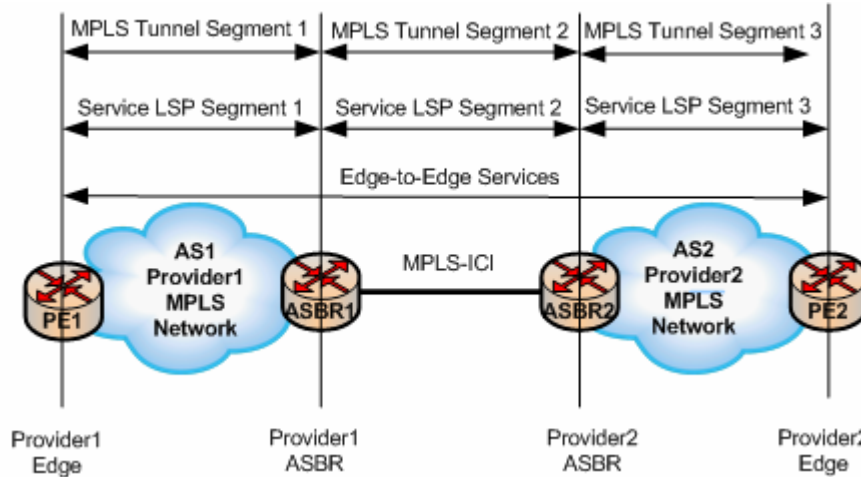
Figure 1 illustrates how BGP Routing and MPLS Signaling (e.g., MP-BGP or RSVP-TE) are used to establish an LSP or LSP segment between two Autonomous System Border Routers (ASBRs) across the MPLS-ICI. It depicts the establishment of tunnel LSP segments across each provider MPLS network and across the MPLS-ICI. The concatenation of these LSP segments forms an edge-to-edge tunnel LSP as shown in Figure 2. Figure 1 also depicts the establishment of service LSP segments (e.g., pseudowire segment or BGP/MPLS IP VPN LSP segment) across each provider MPLS network and across the MPLS-ICI. The concatenation of these service LSP segments forms an edge-edge service LSP as shown in Figure 2. These reference diagrams do not prohibit an edge-to-edge service LSP to be established between Provider Edge (PE)1 and PE2 while being transparent to ASBR1 and ASBR2. Service LSPs, such as PWs for Virtual Private LAN Service (VPLS) and Virtual Private Wire Service (VPWS), and IP VPN labels for Multi-AS backbones Option (b) of MPLS/BGP IP VPNs [RFC4364], are only visible to the ASBR when the hand-off at the ICI is the service label. Service LSPs tunneled over an inter-AS MPLS tunnel are not visible at the ASBR and therefore their processing at the ASBR is out of scope of this specification. These reference models do not mandate that a tunnel LSP segment be established across the MPLS-ICI to establish a service LSP segment across the MPLS-ICI.



**Figure 1: Control Protocols involved in extending MPLS services across an MPLS-ICI interconnecting two provider networks**

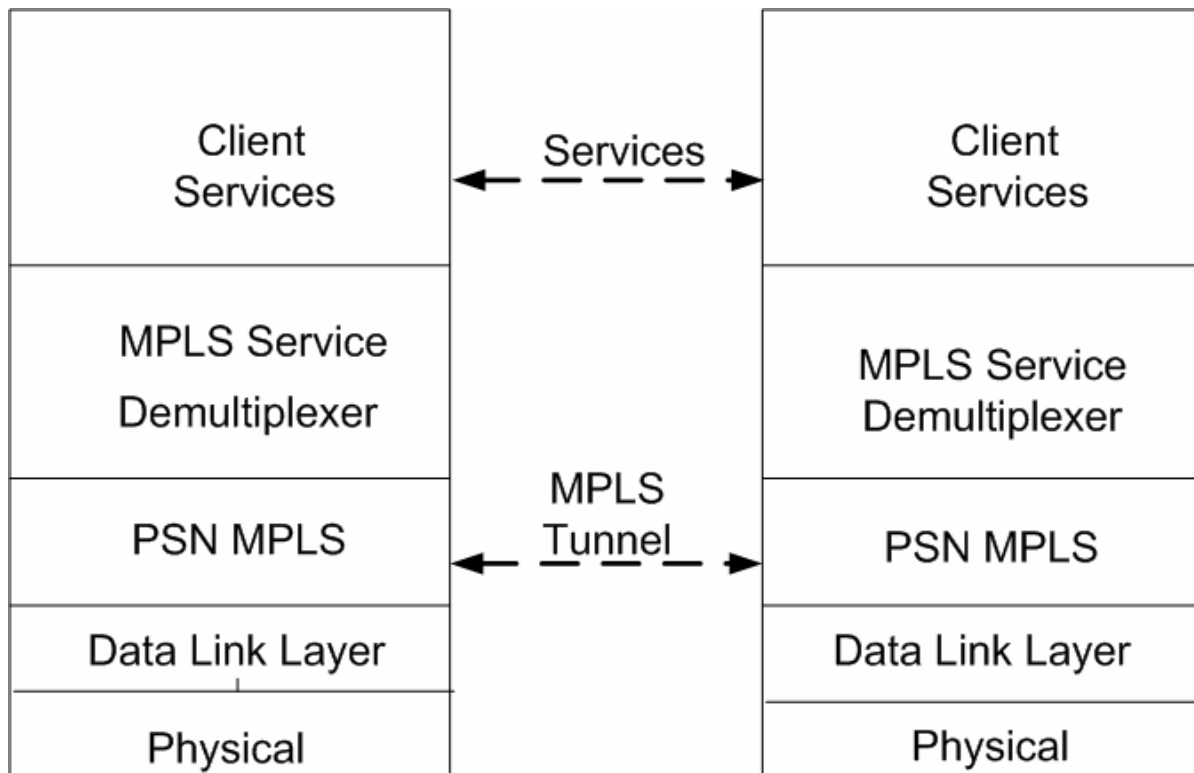
Figure 2 illustrates the edge-to-edge MPLS services and LSPs used to transport these services between two providers. Such services may include (but are not limited to) BGP/MPLS IP VPN, Layer 2 pseudowires (PWs) (e.g., Ethernet PWs), and data trunk tunneling through another Service Provider's network





**Figure 2: Services and LSPs extended across an MPLS-ICI interconnecting two provider networks**

Figure 3 illustrates the protocol stack associated with the MPLS-service transport between two provider networks. Some of the layers in the stack may be transparent to ASBRs, or may be NULL if not used for a given application.



**Figure 3: Protocol Stack**

## 7 Physical layer (Layer 1)

MPLS operates over a variety of physical interfaces including TDM, SONET, Ethernet and optical wavelengths, operating at various speeds (e.g., ranging from T1 to OC192 and beyond). This Technical Specification applies over any physical layer that supports the transport of IP and MPLS packets. An ASBR compliant with this Technical Specification must enable management of Layer 1 interfaces and protocols using the semantics of standard MIBs.

## 8 Data Link Layer (Layer 2)

MPLS operates over a variety of layer 2 technologies including High-Level Data Link Control (HDLC), Point-to-Point (PPP), Packet over SONET, Ethernet and Ethernet VLANs (802.1q/ad), ATM and Frame Relay. This Technical Specification does not mandate or prohibit any specific layer 2 technologies that enable the transport of IP and MPLS packets. An ASBR compliant with this Technical Specification must enable management of layer 2 interfaces and protocols using the semantics of standard MIBs.

## 9 Mechanisms for LSP establishment

This section identifies three mechanisms of LSP establishment across a provider domain boundary:

- all-static configuration
- statically configured and signaled establishment
- dynamic establishment

It specifically identifies mechanisms for Inter-Carrier TE tunnel establishment, Multi-segment PW (MS-PW) establishment, and labeled BGP path establishment for IP VPN and IP routes. These MPLS-services-specific mechanisms, along with associated signaling, routing and management, are discussed in the associated annexes.

### 9.1 All-Static configuration

An ASBR must be able to support static LSPs to establish PWs (PW segments and MS-PWs) and MPLS tunnels across the SP domain (sometimes referred to simply as domain in the rest of the document) boundary. All-Static configuration refers to administratively or manually configuring an LSP or LSP segment that spans the interconnection between two domains' ASBRs. There is no MPLS signaling protocol between these ASBRs, and MPLS labels are administratively or manually assigned. All-Static configuration is labor intensive and introduces challenges in providing for resiliency or reroute around failures, but may be required to satisfy a provider's security and operational requirements. These security requirements may prohibit signaling between domains or require hiding the IP reachability of a PE that is at the far-end of the LSP from the other domain(s). All-Static configuration can be used to establish an LSP that spans the interconnection between two ASBRs and terminates at these ASBRs, or to establish an LSP segment that gets stitched to one or more LSPs in the other domains on either side of the MPLS-ICI.

As stated in sections 5.2 and 5.3 of RFC 4447 [RFC4447], "a (bidirectional) pseudowire consists of a pair of unidirectional LSPs, one in each direction". Section 1 of RFC 4447 states that "packets that are transmitted from one end of the pseudowire to the other are MPLS packets which must be transmitted through an MPLS tunnel. However, if the pseudowire endpoints are immediately adjacent and penultimate hop popping (PHP) behavior is in use, the MPLS tunnel may not be necessary." Therefore,

the procedures for setting up the LSPs and associated resiliency in the case of all-static configuration are treated uniformly in this section.

### 9.1.1 Resiliency

Resiliency of statically-configured LSPs or LSP segments across an MPLS-ICI requires protection against link failure between the same ASBR pair interconnected by the MPLS-ICI or end-to-end LSP protection. End-to-end protection is out of the scope of this document.

Upon the failure of an MPLS-ICI carrying an LSP or LSP segment that is statically configured between two ASBRs, an ASBR should be able to redirect the LSP to a redundant MPLS-ICI. An implementation may support admission control of the LSP over the redundant MPLS-ICI prior to rerouting. Protection against MPLS-ICI failure must be configured on an LSP basis. Where there is no parallel MPLS-ICI between the two ASBRs that can satisfy the constraints of the LSP (e.g., bandwidth) the LSP must be rerouted through other intermediate nodes between the same ASBR pair. Such a redundant path should be configured at the time of the LSP/LSP segment configuration.

#### 9.1.1.1 Protection Models

This Technical Specification describes several protection models that apply to statically configured LSPs, including tunnel LSPs and LSPs that form MPLS pseudowires. These models do not cover layer 1 or layer 2 protection switching. If such protection switching is present, consideration should be given to setting appropriate timer values to facilitate multi-layer switching. Multilayer protection switching, guidance can be derived from section 3.5 of RF 3386 [RFC3386] which states:

*“Multilayer interaction is addressed by having successively higher multiplexing levels operate at a protection / restoration time scale greater than the next lowest layer”.*

If layer1 or layer2 protection is enabled in addition to IP/MPLS protection on an ASBR, the ASBR must be able to apply a method so that each of these lower layers delays failure notification to the next higher layer until it has an opportunity to succeed. Whenever possible, protection switching at the layers underneath the MPLS-ICI should be transparent to the MPLS-ICI.

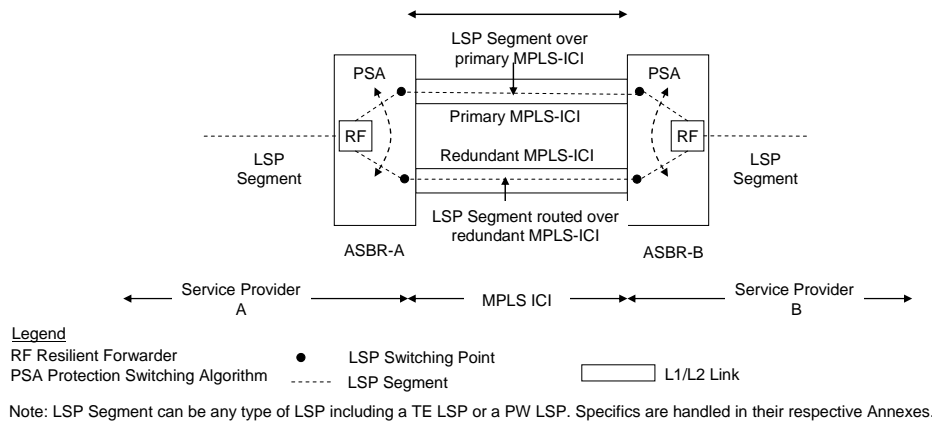
ASBRs compliant with this Technical Specification must support the one-hop MPLS-ICI protection model. ASBRs compliant with this Technical Specification should also support the following protection models: multi-hop LSP protection and multi-hop tunnel bypass described in the following sections.

##### 9.1.1.1.1 One-hop MPLS-ICI protection

One-hop MPLS-ICI protection is illustrated in

Figure 4: One-hop MPLS-ICI Protection. In this case, two ASBRs are interconnected by a pair of MPLS-ICIs. Ideally, these MPLS-ICIs should be diversely routed over layer1 and layer2 networks to reduce fate sharing. The LSP segment is provisioned over a primary MPLS-ICI and a redundant MPLS-ICI. The LSP segment headend shall route the LSP segment over the primary or redundant MPLS-ICI depending on the state of the MPLS-ICI path or the state of the LSP segment. MPLS-ICI failure detection can be based on lower layer mechanisms or IP BFD [IP-BFD]. If the LSP segment has any constraints (e.g., bandwidth), the LSP segment should be admitted over both the primary and redundant MPLS-ICIs (from LSP segment aspect). The LSP segment over the redundant MPLS-ICI should be set up prior to the failure and activated upon failure detection of the primary MPLS-ICI. Upon failure detection and notification of the

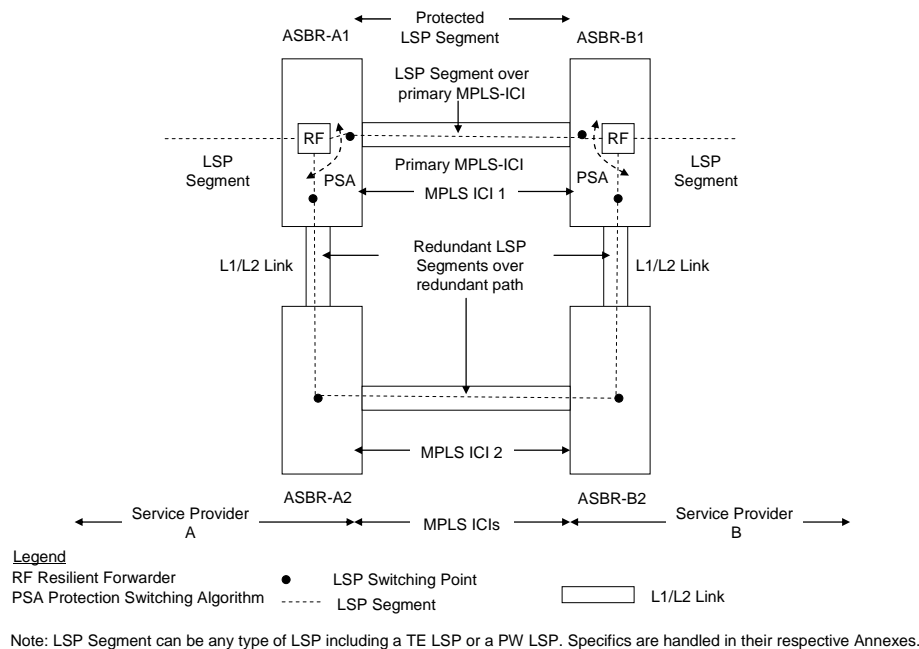
current LSP segment and/or the MPLS ICI over which it is being sent, the LSP packets shall be rerouted over the other path as specified in 9.1.1.3.



**Figure 4: One-hop MPLS-ICI Protection**

#### 9.1.1.1.2 Multi-hop LSP protection over an MPLS-ICI

Multi-hop LSP protection over an MPLS-ICI is illustrated in Figure 5. In this case, an LSP segment is setup between two ASBRs over a direct primary MPLS-ICI interconnecting the two ASBRs. The LSP segment is also set up between the same ASBR pair across multiple hops by configuration, with switching points connecting the segments (e.g., as shown in Figure 5). The multiple segment hops are ASBRs. ASBR-A2 and ASBR-B2 must be configured so that ASBR-A1 and ASBR-B1 use the same LSP label for both the primary MPLS ICI and the multi-hop redundant path. Multi-hop protection may be required when two ASBRs do not have more than one direct MPLS-ICI interconnecting them, or when none of the existing parallel direct MPLS-ICIs can satisfy the LSP constraints, if any. Upon failure detection and notification of the current LSP segment and/or the MPLS ICI over which it is being sent, the LSP packets shall be rerouted over the other path as specified in 9.1.1.3.



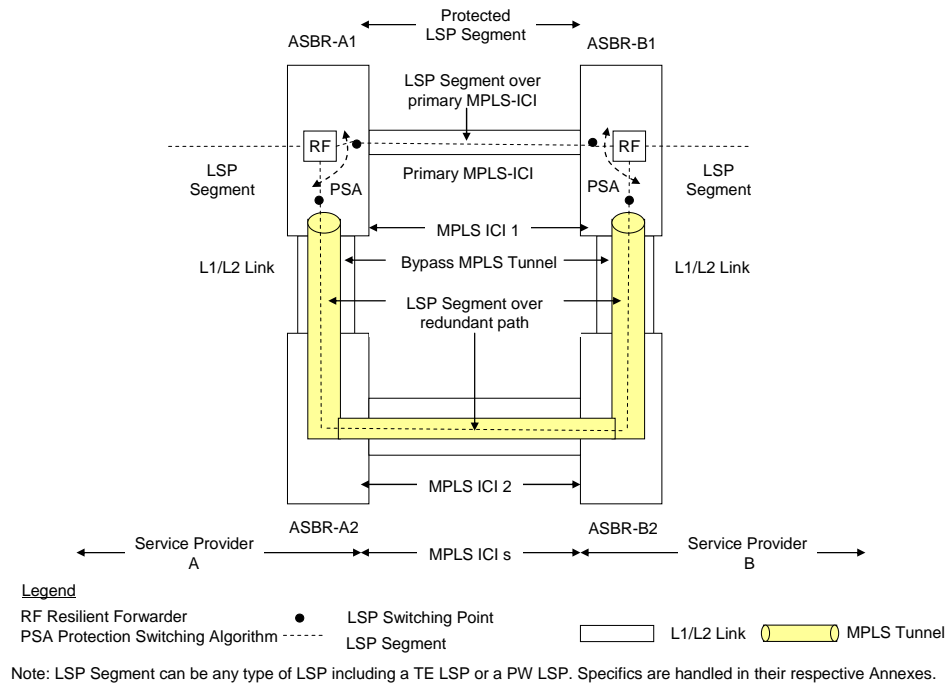
**Figure 5: Multi-Hop LSP protection over an MPLS-ICI**

### 9.1.1.1.3 Multi-hop tunnel bypass protection for the primary MPLS-ICI

Multi-hop tunnel bypass protection for the primary MPLS-ICI is illustrated in Figure 6. An LSP bypass tunnel is configured between two ASBRs to protect one or more MPLS-ICIs interconnecting the same ASBR pair. The LSP bypass tunnel is routed over a layer1 or layer2 path which is diverse from the primary MPLS-ICI and may span one or more other ASBRs but still provide connectivity between the same ASBR pair. The two ASBRs that initiate and terminate the LSP segment are configured to associate that LSP segment with the bypass tunnel for the purpose of protection.

When the head-end of the LSP segment receives an MPLS-ICI failure notification message, it will redirect the LSP packets to the bypass tunnel. The bypass tunnel LSP can carry the protection traffic of one or more LSP segments. The intermediate ASBR nodes only see the bypass tunnel and not the individual LSP segments tunneled through it, as illustrated by the tunnel in Figure 6. These intermediate ASBRs are not required to support individual LSP switching points along the path for the protected LSPs. In addition, the same MPLS label used on the primary MPLS-ICI for a given LSP segment is used for that LSP segment when it is tunneled over the bypass tunnel since the two ASBRs terminating that segment are still directly connected by the bypass tunnel. This capability should be configurable on a per LSP segment basis (i.e., not all LSP segments established over a protected MPLS-ICI are protected).

When associating an LSP segment with a bypass tunnel, that LSP segment should be admitted to the bypass tunnel based on tunnel available resources and attributes and the LSP segment constraints if any. Note that the bypass tunnel must be established before protection is triggered. One MPLS bypass tunnel may protect many LSP segments, but the bypass tunnel need only be established once before LSP segments can be protected by it. Upon failure detection and notification of the current LSP segment caused by the failure of the MPLS-ICI or the bypass tunnel over which it is transported, the LSP packets shall be rerouted over the other path as specified in 9.1.1.3..



**Figure 6: Multi-Hop bypass tunnel (Many:One) protection for the MPLS-ICI**

#### 9.1.1.1.4 End-to-end LSP path protection

For end-to-end LSP path protection, there may be cases where the protection switching function does not reside in the ASBR. In this case, the only role that ASBRs can play in activating protection is failure detection and notification or failure notification relay. End-to end protection may be used simultaneously with other LSP segment protection mechanisms described earlier to obtain a more optimum end-to-end path while effecting fast reroute of the traffic at the point of failure. However it may be necessary to implement appropriate hold-off timers or similar measures for the end-to-end protection instances when operating in conjunction with other forms of protection configured on intermediate segments.

#### 9.1.1.2 Failure Detection and Notification

For all statically configured LSPs, failure detection of the MPLS-ICI can be accomplished by detecting the loss of the physical layer signal, layer 1 OAM, layer 2 OAM, or by utilizing IP Bidirectional Forwarding Detection (IP-BFD) [IP-BFD]. Failure notification for LSP segment protection switching can be based on Layer 1 or Layer 2 OAM, or by utilizing IP-BFD.

ASBRs must support one-hop IP-BFD at minimum as a mechanism for unidirectional and bidirectional MPLS-ICI failure detection. When a failure is detected, local protection may be triggered, if configured, using the one-hop and multi-hop MPLS ICI protection measures described in the previous section. If protection switching is not successful, then an ASBR must send a notification message to the preceding segment of an LSP using either LSP signaling or an MPLS OAM mechanism if present. Pseudowire fault notification messages (e.g. VCCV) are not used to initiate LSP segment protection switching. More details for the pseudowire case are provided in Annex C. Alternatively, end-to-end OAM continuity check or loopback tests may be continuously run to detect failures on the LSP path that could not be locally corrected, but this is out of the scope of this document. MPLS OAM mechanisms for LSPs, other

than those composing MPLS PWs, are described in Section 12 while PW OAM mechanisms are described in Annex C.

### 9.1.1.3 Switchover

The LSP segment MPLS-ICI protection models described in Section 9.1.1 employ unidirectional protection switching for the MPLS-ICI path to an operational redundant LSP-segment path upon failure detection and notification for the MPLS-ICI as described in Section 9.1.1.2. Determining a path operational status may be done using layer1, layer2, or MPLS OAM mechanisms. Referring to Figure 4 through Figure 6, upon receipt of a defect indication for the MPLS-ICI on which traffic is currently being sent (defined in Section 9.1.1.2), the Resilient Forwarding (RF) function in the ASBR at the head-end of the LSP-segment will switch the packet flow from the currently active LSP segment path to the redundant LSP segment path. An ASBR compliant with this Technical Specification shall be able to achieve switchover times within 100 msec, after receipt of the failure notification, using the protection models described in this section.

### 9.1.2 MIBs

When the LSP is a manually configured tunnel and it is not a part of a pseudowire, the ASBR should support the management information model consistent with the LSR MIB as specified in [RFC3813] for static MPLS LSPs, which specifies read-write access for the following groups:

- ifGeneralInformationGroup
- mplsInterfaceGroup
- mplsInSegmentGroup
- mplsOutSegmentGroup
- mplsXCGroup, mplsPerfGroup, mplsLabelStackGroup
- mplsLsrNotificationsGroup

The ASBR should support the management information model consistent with the read-only module compliance of the LSR MIB as specified in [RFC3813] for dynamic MPLS LSPs. An ASBR should provide read-only support for the following groups: mplsInterfaceGroup, mplsInSegmentGroup, mplsOutSegmentGroup, mplsXCGroup, mplsPerfGroup. The ASBR should also support the management information model consistent with mplsLsrNotificationsGroup. An informative description of these groups and their constituent objects is contained in Appendix I.

When the LSP is one end of a pseudowire segment, the cross-connection of two LSPs in one direction of a pseudowire segment across a router (i.e., stitching) should use the information model corresponding to the pwBasicGroup and pwAttachmentGroup compliance groups of [PWMIB].

When the LSP is one end of a pseudowire segment that connects across an MPLS ICI or is nested within an MPLS tunnel, the information model should support full read-write for the following compliance groups of [PWMPLSMIB]: pwMplsGroup and pwMplsOutboundMainGroup. This supports usage of a pseudowire only, or mapping a PW to a TE or non-TE MPLS tunnel.

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An ASBR should support the information model of [PWMPPLSMIB] for setting the PW label EXP bits to a specified value independently of whether there is a Packet Switched Network (PSN) tunnel.

An ASBR should support an information model for configuring the bandwidth required by a PW segment for use by admission control since there is no corresponding object present in [PWMIB].

An informative description of these PW and PWMPPLS groups and their constituent objects is contained in sections I.6 and I.7, respectively.

The following requirements apply to MPLS tunnels and pseudowire LSPs. For one-hop MPLS-ICI LSP protection, there are no standard MIBs and the router should support an information model that allows at least the following attributes:

- Configuration identifying the primary and redundant MPLS-ICIs
- Indication of the up/down state and alarm event notification of the primary and redundant MPLS-ICIs
- Indication of the MPLS-ICI over which LSP packets are currently being sent

For multi-hop MPLS-ICI LSP protection, there are no standard MIBs and the router should support an information model that allows at least the following attributes:

- Configuration identifying the primary MPLS ICI and the first (interface) hop of the redundant multi-hop path
- Indication of the up/down state and alarm event notification of the primary MPLS-ICI and redundant path
- Indication of the path over which LSP packets are currently being sent
- Note that single hop LSP cross-connects must be provisioned at intermediate routers for the redundant LSP

For multi-hop MPLS-ICI LSP protection via a bypass tunnel, there are no standard MIBs and the router should support an information model that allows at least the following attributes:

- Configuration identifying the primary MPLS-ICI and the bypass tunnel providing the redundant path
- Indication of the up/down state and alarm event notification of the primary MPLS-ICI.
- Indication of the up/down state of the bypass tunnel if known from the LSR MIB.
- Indication of the MPLS-ICI over which LSP packets are currently being sent. Note that single hop LSP cross-connects must be provisioned at intermediate routers for the redundant path

Note that the bypass MPLS tunnel would be provisioned using the LSR MIB [RFC3813] as described previously in this section. In addition, the liveness of the bypass tunnel may be verified using MPLS OAM mechanisms described in this document, and accordingly update the bypass tunnel state as up or down.

All static LSP protection schemes should implement a MIB object for protection hold off timer.



## 9.2 Statically Configured and Signaled LSPs: PWs and TE-Tunnels

A statically-configured and signaled LSP is an LSP that is configured administratively or manually at both ASBRs but that is established via dynamic signaling between the ASBRs. As with the all-static LSP, the statically-configured and signaled LSP spanning the MPLS-ICI between two ASBRs, can be administratively or dynamically stitched to other LSPs in either domain. The advantage of the statically-configured and signaled LSPs is that ASBRs, rather than a human administrator or a management system, manage label assignments. The choice of the signaling protocol will depend on the service/application and/or the provider policy. Specifically, there are two service-dependent mechanisms that fit in this category: (1) targeted LDP for MS-PW signaling discussed in Annex C, and (2) TE-tunnels, discussed in Annex D. Resiliency for this LSP establishment mechanism is covered in Annex C and Annex D.

## 9.3 Dynamically Established LSPs: BGP/MPLS IP VPN Multi-AS Backbones Option B, BGP labeled IPv4 paths, and TE-Tunnels

Dynamic LSPs across a domain boundary are those established without any static (manual) configuration at any intermediate point (e.g., ASBR) of the LSP. Implementations that support this Technical Specification shall provide for the dynamic establishment of LSPs via multi-protocol-BGP (MP-BGP) and TE-tunnels using RSVP-TE. The establishment of label-switched paths using MP-BGP for IP VPN and IPv4 routes is discussed in Annex A and Annex B, respectively. TE-tunnel establishment is discussed in Annex D. Resiliency for this LSP establishment mechanism is covered in Annex A, Annex B, and Annex D.

## 9.4 MIBS

The MIBs and management information for dynamically established LSPs in support of pseudowires are described in Annex C. The MIBs and management information for dynamically established TE LSPs are described in Annex D.

# 10 Connection Admission Control

Admission control must be supported at the MPLS-ICI. The admission control function may be located on any system that has visibility of the LSP traffic requirements, available resources at the ICI, and resource management policies for a given carrier's domain. This system, for example, can be the ASBR itself, an element manager, or a bandwidth management system. An ASBR compliant with this Technical Specification must be able to exercise admission control during the setup of an LSP. This does not preclude the use of off-line Connection Admission Control (CAC) tools. However, the specification of off-line CAC tools is outside the scope of this specification.

Connection admission control applies to all mechanisms of LSP establishment and must be based on policies that apply to a traffic profile per Class of Service (CoS) and/or on other administrative policies. It must be possible to enable/disable admission control based on bandwidth constraints at the MPLS-ICI level. Other policies that affect admission control must be configurable and will vary per MPLS service. Admission control specific to MPLS services within the scope of this document is discussed within the respective annexes. ASBRs compliant with this Technical Specification must support the following generic capability for statically configured LSPs (all-static and statically-configured but signaled) with bandwidth and CoS constraints:

When an LSP with bandwidth and CoS requirements is configured between two ASBRs compliant with this Technical Specification, the ASBRs must be capable of applying admission control to the LSP over one or more TE-tunnels or MPLS-ICIs between the ASBRs to accommodate the bandwidth requirement

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of the LSP per CoS. As a result, if the LSP is admitted, it is bound to the chosen MPLS-ICI or TE tunnel. If the LSP cannot be admitted because of insufficient resources, the ASBR should keep the LSP down and send a trap to a network management element notifying it of the failure. If the LSP being set up is a segment of an end-to-end LSP that extends into either domain on either side of the MPLS-ICI, the end-to-end LSP setup must fail when the LSP segment setup fails.

## 11 Forwarding

### 11.1 Traffic management

A class of service (CoS) is an aggregate for all traffic that shares similar Quality of Service (QoS) objectives and similar characteristics. We use the term CoS in this document interchangeably with service class, in the same manner that has been used in [RFC4594]. The Diffserv architecture [RFC2475], MPLS Diffserv extensions [RFC3270] and Diffserv-aware MPLS traffic engineering extensions [RFC4124] defined by the IETF provide the architecture for supporting the transport and signaling of multiple classes of service [RFC4124][RFC3270]. However, they have not yet addressed most of the issues that arise when establishing inter-domain inter-provider MPLS connections.

In the forwarding plane, a CoS is associated with a Per Hop Behavior (PHB). PHBs as defined in [RFC2475] “are implemented in nodes by means of some buffer management and packet scheduling mechanisms”. In MPLS networks that support Diffserv, the PHB that a packet should receive at a node can be inferred from the MPLS label and EXP bits (L-LSPs) or from the EXP field value (E-LSPs) [RFC3270]. This document addresses E-LSPs only. For IP packets, as in Multi-AS Backbones Option A, the PHB that a packet should receive is identified by the Diffserv Code Point (DSCP) [RFC2475].

#### 11.1.1 Classes of Service and EXP Mappings in the Forwarding Plane

Each SP implements a set of classes of service within its network. SPs may use different EXP-field values to correspond to similar classes of service. SPs may also implement different classes of service in their respective networks. In the data path, the CoS (for E-LSPs) and drop precedence identifiers are encoded in the EXP field of the MPLS header of an MPLS packet. Thus, in the data path, the EXP field value used in one provider network may need to be mapped to another EXP field value used in the other provider network for packets to receive the same treatment across both networks or to perform CoS and drop precedence mappings. When providers implement different classes of service, they may decide to also map one CoS in one network to another CoS in the other network. This CoS mapping is often defined in a bilateral agreement between providers. In the forwarding plane, the agreed-on CoS mapping translates to a mapping between corresponding EXP field values whereby each EXP field value identifies a PHB. The desired externally observable behavior is that packets with a given EXP field value are directed to a particular queue and drop precedence and they leave the ASBR with a particular EXP field value that may possibly be different from the one received at the ASBR.

In order to support inter-provider operations as described earlier, ASBRs that support this Technical Specification must enable the configuration of EXP field value mapping between incoming EXP values and outgoing EXP values. A provider must be able to apply an EXP map at the ASBR ingress for traffic received over an MPLS-ICI and at the ASBR egress for traffic sent over an MPLS-ICI. If the mapping is applied to the egress of an interface, an ASBR performs the mapping and (re)writes the EXP field value in the outer label header and possibly EXP value in the inner label header (based on configuration) with the outgoing EXP field value prior to transmitting the packet to the next ASBR over the MPLS-ICI. If the mapping is applied at ingress, an ASBR performs the mapping when it receives the packet and (re)writes the EXP field value in the outer label header, and possibly inner label header, with the outgoing EXP field

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value when transmitting the packet within its network. Appendix II illustrates via examples how EXP mapping could be configured at ASBRs interconnected by an MPLS-ICI. ASBRs should be capable of also programming the EXP to PHB map dynamically based on signaling a Diffserv-MPLS TE tunnel as defined in [RFC3270] section 5.2.1. When performing label popping or imposition, equipment must support the pipe model, short-pipe model and uniform models as defined in [RFC3270]. The model applied must be configurable per LSP and per interface.

An ASBR compliant with this Technical Specification must be able to associate an EXP field value with a queue and a buffer management profile associated with that queue before or after the EXP field mapping. A queue is served by a scheduling structure. A buffer management profile applied to a queue is used to decide whether to en-queue the packet or drop it based on the state of the queue, the length of the packet and the color of the packet [RFC2475].

### 11.1.2 Traffic Policing

A downstream provider may choose to enforce the traffic contract of an LSP or an aggregate of LSPs by policing traffic forwarded by the upstream provider. ASBRs compliant with this Technical Specification must be able to apply a policer to traffic that belongs to an LSP, a group of LSPs, and/or a CoS identified by an EXP field value on an MPLS-ICI. It must also be able to police traffic per IP VPN Attachment Circuit (AC) per IP Precedence/DSCP (as it applies to multi-AS Backbone Option A). The LSP traffic contract itself can be derived from signaling and/or administrative configuration.

For policers, ASBRs must implement the one-rate-three-color-marker [RFC2697] and the two-rate-three-color-marker [RFC2698] metering algorithms in both color-blind and color-aware modes. The actions taken by the policers can be to pass the packet, drop the packet, or to increase its drop precedence. Marking drop precedence is done in the EXP field value of the MPLS header.

### 11.1.3 Traffic Shaping

If a downstream node polices traffic per LSP, group of LSPs, and/or EXP, or per Attachment Circuit (AC) and/or DSCP/Precedence, an upstream node that sends the traffic over the MPLS-ICI should be able to shape traffic to the traffic profile used in the metering and policing action. An ASBR compliant with this Technical Specification must support the following:

- A single-rate shaper per EXP and per interface
- A single rate shaper per IP VPN attachment circuit and/or IP Precedence/DSCP (as it applies to Multi-AS Backbones Option A)

An ASBR should support the following:

- A single-rate shaper per EXP per LSP or group of LSPs
- Dual-rate shapers

The ASBRs must enable or disable shaping on queues by configuration.

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## 11.2 Path MTU Handling and Fragmentation

Maximum Transmission Unit (MTU) is the largest packet size that an interface can transmit without fragmentation. MTU is measured in Bytes (B) and it is media-dependent. It represents the largest payload that can be carried in a Layer-2 frame, and it does not include Layer-2 headers or trailers.

While MTU is defined with respect to interfaces, the end-to-end packet transport depends on the Path MTU (PMTU), which is defined as the smallest MTU along the packet path from the source to the destination. The packet sender is usually concerned with the PMTU, rather than the MTU of the attached link, because it is the PMTU that determines whether the packet will be delivered to the receiver and whether it will be fragmented in the process.

The PMTU is determined by the Path MTU Discovery (PMTUD) process, which is defined in [RFC1191] for IPv4 (PMTUDv4). PMTUDv4 [RFC1191] operates as follows:

- IP sender sets the “Don’t Fragment” (DF) flag in the IP header.
- IP sender sends a datagram corresponding to the MTU of the first hop.
- When an ASBR is encountered with the next-hop link MTU too small to send the packet without fragmentation, it discards the packet (because DF is set) and sends to the IP sender an Internet Control Message Protocol (ICMP) message “Destination Unreachable” with a code meaning “Fragmentation Needed and DF Set.”
- IP sender receives the ICMP message and sends a smaller packet.
- The process is repeated until a packet is small enough to get through without generating an ICMP message.
- This smallest packet becomes the PMTU used by the IP sender.

It should be noted that PMTUD could potentially be used to create a Denial of Service (DoS) attack on the control plane of the ASBRs that generate ICMP messages in response to packets that exceed the outgoing interface MTU and that cannot be fragmented.

### 11.2.1 Label stack depth

Considering that MPLS operates between Layer 2 and Layer 3, as far as Layer 2 is concerned, the MPLS label stack represents a part of the payload. Therefore, the MPLS label stack may change the status of the Layer 2 payload from “not exceeding the MTU size” to “exceeding the MTU size.” MTU scenarios for MPLS-ICI may further depend on whether a new header is provided for a tunnel on the MPLS-ICI and whether this header represents an additional layer in the label stack. The following tunnel cases may arise:

- Contiguous tunnel – new header on the MPLS-ICI, but the label stack depth does not change.
- Stitched tunnel – new header on the MPLS-ICI, but the stack depth does not change.
- Nested tunnel – new header on the MPLS-ICI, the stack depth increases; and labeled packets may exceed the MTU of the MPLS-ICI as a result of imposing the additional label.

### 11.2.2 Fitting packet sizes to PMTU

Different domains may implement different approaches to fitting packet sizes to the PMTUs, including the following:

1. sending packets that correspond to the minimum MTU (576 Bytes for IPv4 networks)
2. fragmenting packets (in any node along the path for IPv4)
3. PMTUD process to determine the PMTU
4. using the RSVP-TE procedures defined in [RFC3209] Section 2.6 (Path MTU) when signaling an RSVP-TE LSP
5. using LSP trace [RFC4379] and the downstream mapping object included in echo reply

Approach-1 is wasteful of bandwidth if the PMTU is larger than 576 Bytes. Approach-2 is wasteful of network processing resources and creates delay that may be unacceptable for real time applications that send large packets, e.g., video conferencing. Approach-3 is most efficient compared to approach-1 and approach-2 from bandwidth and processing viewpoint. The only drawback of option-3 is that it could result in attacks on a router control plane if that router is not properly protected. Approach-4 is best suited when signaling an RSVP-TE LSP as this approach does not require additional processing on routers processing the RSVP-TE messages. Approach-5 could be useful when using LSP ping trace-mode for path trace. In addition, LSP ping trace-mode could be specifically used for this processing. However, as stated in section 12, LSP trace could be construed as a potential DoS attack on routers.

### 11.2.3 MTU Requirements for MPLS-ICI

Following are requirements that pertain to MTU handling on the MPLS-ICI:

- An ASBR should support PMTU Discovery for IPv4.
- The PMTU Discovery process for MPLS tunnels must follow the process described in [RFC3032], Section 3.6.
- An ASBR compliant with this Technical Specification must be able to create, receive and pass "Destination Unreachable" ICMP messages with a code meaning "Fragmentation Needed and DF Set".
- SPs must be able to configure an option for silently dropping (i.e., without generating ICMP messages) packets whose lengths exceed the outgoing interface MTU.
- An ASBR compliant with this Technical Specification must be able to detect whether the payload of an MPLS packet is an IP packet or not in order to perform the right fragmentation procedure if fragmentation is permitted.
- An ASBR compliant with this Technical Specification must resist Denial of Service (DoS) attacks that use ICMP messages. In particular, an ASBR must be able to rate limit the number of generated ICMP messages when the packet forwarder at the ASBR sends packets to an interface with a size that exceeds the interface MTU.

- An ASBR compliant with this Technical Specification should support a configuration parameter called “Maximum Initially Labeled IP Datagram Size” as described in RFC 3032, Section 3.2.
- An ASBR compliant with this Technical Specification must implement the procedures defined in RFC 3032, Section 3.4, for processing IPv4 datagrams that are too big.
- An ASBR compliant with this Technical Specification must support the MTU processing as required by Section 2.6 of RFC 3209 [RFC3209]
- An ASBR compliant with this Technical Specification must support MTU processing as required for LSP ping trace-mode [RFC4379]. In particular, the MTU returned by an ASBR in the echo reply must be the minimum MTU of the path to the next hop (and associated label map) returned by the ASBR.

[RFC4623] defines generic PW encapsulation and procedures that allow a PW packet fragmentation before insertion of the PW into the PSN at the ingress PE and the reassembly of the PW packet fragments at the egress PE. Similar procedures can be applied at an ASBR if the ASBR is a switching point (S-PE) of the PW (c.f., see [MS-PW\_Requirements] for S-PE). An ASBR compliant with this Technical Specification should support [RFC4623]. In this case, an ASBR may segment and/or reassemble a PW packet. The choice of the outgoing MTU may be discovered or configured. In the discovery model, an ASBR may be able to use the path MTU discovery mechanisms to discover the path MTU of a PW segment that extends between two ASBRs for instance. If the PW is an IP PW, procedures described in [RFC3032] and discussed earlier in this section are applicable. An ASBR that complies with this Technical Specification must have the capability of turning off PW fragmentation if that capability exists and should provide the capability for allowing fragmentation/reassembly selectively for certain pseudowires but not others.

### 11.3 Load Balancing

An ASBR compliant with this Technical Specification must be able to load balance MPLS traffic across equal cost providers’ paths, including interconnects. The ASBR must enable the operator to define the number of topmost consecutive labels in the label stack used in load-balancing MPLS labeled packets. A minimum of 2 topmost labels must be supported. An ASBR must be capable of enabling/disabling load balancing and defining load balancing criteria when enabled. Implementations must ensure that packets from the same microflow do not get load-balanced across different paths when load balancing is enabled. For instance, packets from the same PW in a given direction must follow the same path. If load balancing involves MPLS labels only, this can be ensured. If the load balancing criterion includes IP information, the forwarding path must be able to detect whether the payload of an MPLS packet is IP [RFC4385] and apply load balancing only when applicable.

It should be noted that Equal Cost Multi-Path (ECMP) at intermediate nodes as a matter of packet forwarding applies to non-reserved traffic, that is, traffic whose path is not reserved. ECMP also introduces challenges in tracing data paths and conducting performance measurements across all possible paths. These challenges are operational in nature.

### 11.4 Time To Live (TTL) Processing

[RFC3443] supports multiple models for TTL processing in hierarchical MPLS networks: (1) uniform, (2) short pipe with and without Penultimate Hop Popping (PHP), and (3) pipe. An ASBR in conformance to this Technical Specification must support all these models. Specifically, the pipe model and short-pipe

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model may be most useful for providers who like to hide the length of the path across their network when tunneling LSPs and/or IP packets through MPLS tunnels in their networks.

## 11.5 MIBs

The configuration and management of Diffserv-related traffic management components are specified in [RFC3289]. An informative description of how these components can work together is given in [RFC3290]. The Diffserv MIB defines a means to define a sequence of components on either the ingress or egress direction of an interface. The MIB defines some of these components as well as provides a generic row pointer mechanism to point to other MIBs that define the configuration and management of other components (e.g., an MPLS label and EXP bit classifier Table). The generic components defined in [RFC3289] are: classification, metering, action, and queuing. The ASBR should support the management information model consistent with the following groups with read-write access as specified in [RFC3289]:

- diffServMIBDataPathGroup
- diffServMIBClfrGroup
- diffServMIBClfrElementGroup
- diffServMIBMultiFieldClfrGroup
- diffServMIBActionGroup
- diffServMIBAlgDropGroup
- diffServMIBQGroup
- diffServMIBSchedulerGroup
- diffServMIBMaxRateGroup
- diffServMIBMinRateGroup
- diffServMIBCounterGroup
- diffServMIBMeterGroup
- diffServMIBTBParamGroup
- diffServMIBDscpMarkActGroup
- diffServMIBRandomDropGroup

An informative description of these objects, their indexing, content, suggested usage and context is contained in Appendix I.

Since diffServMultiFieldClfrTable in [RFC3289] does not cover MPLS labels and EXP bits, the ASBR should support an enterprise MIB or equivalent information element that contains classification

parameters that match certain EXP bits and LSP label values for an MPLS packet in the ingress and egress directions of an interface as specified in section 11.1.1. The DiffServClfrEntry associated with these MPLS labels and EXP parameters would then specify the next Diffserv component (e.g., a meter).

Since diffServDscpMarkActTable only (re)marks the DSCP value in the IP packet header, there is a need for an Enterprise MIB or equivalent information model that contains the value of the EXP markings to be applied in the MPLS packet header.

## 12 OAM

The MPLS OAM functionality can be split into two types: (1) ‘always-on’ defect detection and handling and (2) ‘on-demand’ diagnostics. OAM is vital to network operations and impacts provider Operational Support Systems (OSS). Due to its importance, ASBRs compliant with this Technical Specification must support both MPLS OAM functions: defect detection and handling as well as diagnostics. The defect detection and handling function needs to be as simple as possible to minimize the processing cost. The diagnostics tools must include functions such as echo request/reply and path trace.

This section discusses OAM capabilities at an MPLS-ICI related to all of the MPLS services described in the annexes. OAM processing requirements in this section only apply to OAM packets that are destined to ASBRs interconnected by an MPLS-ICI. The focus of this section is OAM capabilities that apply to LSP segments established across the MPLS-ICI and to other segments of the same LSP that extend beyond the MPLS-ICI. End-to-end OAM is out of the scope of this document. The LSP segment established across the MPLS-ICI is a segment of an LSP that extends from PE to PE across ASes. This LSP segment is dynamically or statically established and stitched.

### 12.1 Connection Verification: “Always On” Defect Detection and handling

ASBRs compliant with this Technical Specification must support Bidirectional Forwarding Detection (BFD) [MPLS-BFD] in asynchronous mode. They must support the following for BFD:

- Timer Parameters:
  - time interval between successively transmitted protocol messages per LSP.
  - minimum receive interval for protocol messages per LSP.
  - failure detection criterion in terms of the number of successive messages that must be lost to trigger the declaration of an LSP down. This must be configurable per LSP.
- Failure notification
  - An ASBR should be able to send an SNMP trap upon LSP failure detection.
  - An ASBR must notify all client protocols that depend on the liveness of the LSP being monitored when that LSP fails.

An ASBR acting as an endpoint for an LSP or LSP segment must be able to avert OAM-related DoS attacks by:

- Dropping the BFD protocol messages received on an LSP if the protocol is not enabled for that LSP.



- Policing BFD protocol messages to enforce the message rate configured for all LSPs on an MPLS-ICI. Per-LSP policing is optional.

It should be noted that policing BFD messages used for liveness check may result in a false failure detection. Thus, the policing parameters must be set so that “legitimate” messages used for liveness check are not impacted by policing unless they exceed their allocated rate.

ASBRs compliant with this Technical Specification must allow enabling an MPLS BFD session per LSP. When the ASBRs are transit points for an LSP, BFD running between the LSP endpoints is outside of the scope of this document. An LSP extended between ASBRs can be stitched to other LSP segments to form an end-end LSP. In that case, a BFD session that runs end-to-end between the LSP endpoints is transparent to the ASBRs. However, ASBRs must support the co-existence of an end-to-end BFD session and a BFD session between the ASBRs for an LSP segment. MPLS packets carrying the BFD messages corresponding to the ASBR BFD sessions must have TTL set to 1 to force these messages to be processed at the ASBRs rather than be switched across.

ASBRs should support the authentication option for multi-hop and single-hop BFD sessions between two ASBRs as discussed in [BFD-Base]. Where applicable, authentication must be enabled by configuration and the same key or password should be sharable for all sessions between the same ASBR pair using the same authentication method.

ASBRs compliant with this Technical Specification must support the bootstrapping method via MPLS ping for exchanging Your Discriminator and My Discriminator values used in BFD control messages. An ASBR compliant with this Technical Specification should also support the configuration of the local Discriminator for a BFD session (My Discriminator value in the BFD messages the ASBR sends to the peer ASBR at the other end, and Your Discriminator value in the BFD messages it receives for the session), and the configuration of the peer ASBR Discriminator for a BFD session (My Discriminator value in the BFD messages the ASBR receives from the peer ASBR at the other end, and Your Discriminator value in the BFD messages it sends for the session).

## 12.2 On Demand Diagnostics

ASBRs compliant with this Technical Specification must support LSP ping [RFC4379] in both ping and trace modes to verify unidirectional connectivity and perform path tracing of MPLS label switched paths on MPLS-ICI segments, respectively. Equipment compliant with this Technical Specification should also support BFD in echo mode [MPLS-BFD] for performing loopback tests.

ASBRs compliant with this Technical Specification must support LSP ping [RFC4379] in ping mode for checking the liveness of the LSP and the data plane state against the control plane state for that LSP. This applies to LSPs with endpoints on the ASBRs at either end of an MPLS-ICI. That LSP may be a segment of an end-to-end LSP that extends beyond the MPLS-ICI. LSP ping messages for LSPs that do not terminate on an ASBR MPLS-ICI transit the MPLS-ICI. An ASBR that is compliant with this Technical Specification and supports a service described in the annexes must support the associated LSP ping Forwarding Equivalence Class (FEC) sub-Type Length Value (sub-TLV) and their stackings as indicated in the following table:

**Table 1 - LSP ping FEC and sub-TLVs**

<u>Sub-type</u>	<u>Length</u>	<u>Field value</u>	<u>Application</u>
3	20	RSVP IPv4 LSP	TE-Tunnels
6	13	VPN IPv4 prefix	IP VPN
8	14	L2 VPN endpoint	Pseudowires
10	14	"FEC 128" Pseudowire	Pseudowires
11	16	"FEC 129" Pseudowire	Pseudowires
12	5	BGP labeled IPv4 prefix	Labeled IPv4 routes

LSP ping must support the following reply modes specified in [RFC4379]:

**Table 2 - LSP ping reply modes**

<u>Value</u>	<u>Meaning</u>
1	Do not reply
2	Reply via IPv4 UDP packet
3	Reply via applications level control channel

LSP ping replies could include replies with the router alert option. Such replies will cause every router on the path of the LSP to process the LSP ping messages. In order to prevent ping replies originated in one provider domain to be processed on every router in another provider domain on the path to the destination, the reply mode should not include the router alert option. An ASBR must be configurable to drop or rate-limit received echo reply packets with the router alert option to avert overloading or attacking the ASBR control plane and that of other routers within the ASBR AS. Another way for a network to avoid DoS attacks is to transparently pass the packets with the router alert option. However, this will also prevent processing other packets with the router alert option.

The following must also be supported for LSP ping in ping mode:

- Timer parameters:
  - time interval between successive echo requests per LSP or globally when initiating an LSP ping test.
  - failure detection criterion in terms of the number of successively missed echo request replies that triggers the declaration of an LSP down. This must be configurable per LSP.
- Failure notification
  - An ASBR should be able to send an SNMP trap upon LSP failure detection.

- An ASBR should notify all client protocols that depend on the liveness of the LSP being monitored when that LSP fails.

An ASBR acting as an endpoint for an LSP must be able to avert OAM-related DoS attacks by:

- Dropping the LSP ping messages received on an LSP if the protocol is not enabled for that LSP.
- Policing LSP ping echo requests to enforce the message rate configured for all LSPs. Per-LSP policing is optional.

Trace mode capability of LSP ping can be used for fault isolation. LSP path tracing enables the identification of the path(s) traversed by an LSP and hop-by-hop fault localization. ASBRs compliant with this Technical Specification must provide the capability to rate limit or drop LSP tracing messages arriving at an ASBR from another provider to be processed at the ASBR. When an ASBR drops an LSP ping message it will disrupt the end-to-end path trace. An ASBR should also support the option to respond at the domain boundary without including a downstream label map.

When responding to path trace messages from another provider ASBR, an ASBR must be configurable to respond with either of the following two options: (1) respond without a downstream label map to the next hop, (2) respond with a downstream label map to the next hop. Protection of confidential information when the path trace message is targeted to a router behind the ASBR within an AS is still an issue to be addressed. The detection of the LSP ping messages at an ASBR may require deep packet inspection and may not be feasible in certain cases. If a router inside an AS has knowledge that the LSP being traced is a cross-AS LSP it may drop the LSP ping echo request or respond to the LSP ping echo request without the downstream label map.

### 12.3 MIBS

The ASBR should support a management information model that provides for the configuration and management of BFD in a way that is consistent with [BFD MIB] compliance for the following groups:

- bfdSessionGroup,
- bfdSessionPerfGroup,
- bfdSessionPerfHCGroup,
- bfdNotificationGroup

An informative description of these objects, their indexing, content, suggested usage and context is contained in Appendix I.

The information model in the BFD MIB contains objects that can be configured to cause an ASBR to generate a significant number of packets, potentially creating a DoS attack on other routers or networks. Therefore, when the MIB is used the ASBR shall support only SNMPv3 for configuration of the BFD MIB. When read-only operations are performed SNMPv2 is sufficient. When other network management protocols are used, a mechanism with a comparable level of security should be used.

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## 13 Security and Confidentiality

This section discusses security capabilities that are important at the MPLS-ICI, and at devices (including ASBRs) which support the ICI. Security threats, capabilities, and practices that occur or may be deployed elsewhere in the network are outside of the scope of this section. More network-wide or general discussions of threats, capabilities, and practices can be found in [RFC4111] [RFC3871] [RFC4778]. The security capabilities stated in this section should be considered as complementary to security considerations addressed in the individual protocol specifications and/or security frameworks. This section discusses threats and capabilities at the Inter-Carrier Interconnect related to all of the use cases described in the annexes later in this document.

Security vulnerabilities and exposures may be propagated across multiple networks because of security vulnerabilities arising in one peer's network. Threats to security originate from accidental, administrative and intentional sources. Intentional threats include events such as spoofing and DoS attacks.

The level and nature of threats, as well as security and availability requirements, may vary over time and from network to network. This section therefore discusses capabilities that need to be available in ASBRs deployed for support of the MPLS-ICI. Whether any particular capability is used in any one specific instance of the ICI is up to the service providers managing the ASBRs offering/using the ICI services.

### 13.1 Control Plane Protection

This section discusses capabilities for control plane protection, including protection of routing, signaling, and OAM.

#### 13.1.1 Authentication of Signaling Sessions

ASBRs compliant with this Technical Specification must support MD5 authentication for all TCP-based protocols within the scope of the MPLS-ICI (i.e., LDP signaling, and BGP routing) and MD5 authentication for the RSVP-TE Integrity object to interoperate with current practices.

An ASBR compliant with this Technical Specification should be able to support exchange of all signaling and routing (LDP, RSVP-TE, and BGP) protocol messages over a single IPSec tunnel in tunnel or transport mode with authentication but with NULL encryption, between the peering ASBRs. IPSec, if supported, must be supported with HMAC-MD-5 and optionally SHA-1. It is expected that authentication algorithms will evolve over time and support can be updated as needed. An ASBR must support key configuration for an IPSec session.

OAM operations across the MPLS-ICI could also be the source of security threats on the provider infrastructure as well as the service offered over the MPLS-ICI. A large volume of OAM messages could overwhelm the processing capabilities of an ASBR if the ASBR is not properly protected. Maliciously-generated OAM messages could also be used to bring down an otherwise healthy service (e.g., MPLS pseudowire), and therefore affect service security. LSP ping [RFC4379] does not support authentication today and that support should be subject for future considerations. BFD [BFD-Base] however does have support for carrying an authentication object. It also supports TTL processing as an anti-replay measure. Implementations conformant to this MPLS-ICI should support BFD authentication using MD-5 and must support the procedures for TTL processing.

### **13.1.2 Protection against DoS attacks in the Control Plane**

An ASBR compliant with this Technical Specification must provide the ability to filter signaling, routing, and OAM packets destined for itself, and must provide the ability to rate limit such packets. Packet filters should be capable of being separately applied per interface, and should have minimal or no performance impact on the ASBR. For example, this capability allows an SP to filter and rate-limit signaling, routing, and OAM messages that can be sent by a peer SP to an associated traffic profile.

In the presence of a control plane DoS attack against an ASBR, the ASBR should guarantee sufficient resources to allow network operators to execute network management commands to take corrective action, such as turning on additional filters or disconnecting an interface which is under attack. DoS attacks on the control plane should not adversely affect data plane performance.

ASBRs which support BGP must support the ability to limit the number of BGP routes received from any particular peer. Furthermore, in the case of IP VPN, an ASBR must be able to limit the number of routes learned from a BGP peer per IP VPN. In the case that an ASBR has multiple BGP peers, it should be possible for the limit to vary between peers.

### **13.1.3 Protection against Malformed Packets**

ASBRs compliant with this Technical Specification should be robust in the presence of malformed protocol packets. For example, malformed routing, signaling, and OAM packets should be treated in accordance with the relevant protocol specification.

### **13.1.4 Ability to Enable/Disable Specific Protocols**

ASBRs compliant with this Technical Specification must allow an administrator to enable or disable a protocol (by default, a protocol is disabled unless administratively enabled) on a per interface basis. ASBRs must be able to drop any signaling or routing protocol messages when these messages are to be processed by the ASBR but the corresponding protocol is not enabled on the interface on which the messages were received. This dropping should not adversely affect data plane or control plane performance.

### **13.1.5 Protection against Incorrect Cross Connection**

ASBRs compliant with this Technical Specification must support LSP ping [RFC4379]. This may be used to verify end to end connectivity for the LSP (e.g., PW, TE Tunnel, VPN LSP, etc), and to verify PE to PE connectivity for L3 VPN services.

ASBRs and Route Reflectors (RRs) which support BGP operation must allow a means to restrict which Route Target attributes are sent to and accepted from a BGP peer across an ICI. ASBRs and RRs should also be able to inform the peer regarding which Route Target attributes it will accept from the peer [BGP-ORF]. This is due to the fact that an incorrect Route Target sent by the peer can result in incorrect cross-connection of VPNs. Also, sending inappropriate route targets to a peer may disclose confidential information. Further security considerations for inter-provider BGP/MPLS IP VPN operations are discussed in Annex A.

### **13.1.6 Protection Against Spoofed Updates and Route Advertisements**

ASBRs compliant with this Technical Specification must support signaling and routing authentication per section 13.1.1. They must also support filtering of routes received via a BGP peering session by applying

policies that include one or more of the following: AS path, BGP next hop, standard community and/or extended community.

### **13.1.7 Protection of Confidential Information**

ASBRs compliant with this Technical Specification should provide the ability to identify and prohibit messages that can reveal confidential information about network operation (e.g., performance OAM messages, LSP Traceroute messages). Service Providers must have the flexibility of handling these messages at the ASBR. For example, ASBRs supporting LSP Traceroute may limit addresses to which replies can be sent. Note that this capability should be used with care. For example, if a service provider chooses to prohibit the exchange of LSP ping messages at the MPLS-ICI, it may make it more difficult to debug incorrect cross-connection of LSPs or other problems.

A provider may decide to progress these messages if they are incoming from a trusted provider and are targeted to specific agreed-on addresses. Another provider may decide to traffic police, reject or apply policies to these messages. Solutions must enable providers to control the information that is relayed to another provider about the path that an LSP takes. For example, in RSVP-TE record route object or LSP ping trace, a provider must be able to control the information contained in corresponding messages when sent to another provider.

## **13.2 Data Plane Protection**

### **13.2.1 Protection against DoS in the Data Plane**

This is provided via traffic policing as described in Section 11.1.2.

### **13.2.2 Protection against Label Spoofing**

ASBRs compliant with this Technical Specification must be able to verify that a label received across an MPLS-ICI was actually assigned to an LSP arriving from the provider across that MPLS-ICI. If the verification fails, the corresponding MPLS packet must be dropped. This verification can be applied to the top label only. The top label is the received top label and every label that is exposed by label popping to be used for forwarding decisions.

ASBRs compliant with this Technical Specification must provide the capability of dropping MPLS-labeled packets if all labels in the MPLS label stack are not processed at the ASBR. The presence of non-processed MPLS headers is detected if the S-bit in the last processed MPLS header is set to 0. This behavior must be configurable on a per interface basis. Enabling this behavior on an interface prevents some applications across that interface. However, when enabled, it provides an SP the capability of guaranteeing that every label that enters its domain from another SP was actually assigned to that SP and avoids a potential security attack on a service within its domain.

## **Annex A IP VPN**

This annex covers the IP VPN MPLS service based on BGP/MPLS IP VPNs as defined in [RFC4364]. In particular, this annex discusses option A and option B of the ‘Multi-AS Backbones’ section of [RFC4364].

There are cases where a single IP VPN provider may not have the footprint necessary to serve all sites of an enterprise IP VPN customer. Thus, providers find themselves forming partnerships with other providers, the nature of which depends for example on mutual business benefit or on an enterprise demanding that a certain service provider be included in a specific solution.

Multi-AS Backbones Option A, ‘VRF-to-VRF connections at the AS (Autonomous System) border routers’, is described in section 10 item (a) of [RFC4364]. Multi-AS Backbones option A requires that unlabeled IPv4 unicast addresses be exchanged between the ASBRs and that forwarding be based on information in the IP header of the received IP datagram. Customer DSCP/IP precedence preservation is especially important when the two interconnected providers’ networks use different DSCP values to indicate the same Per-Hop-Behaviors (PHBs) or use different sets of PHBs.

Multi-AS Backbones Option B, ‘EBGP redistribution of labelled VPN-IPv4 routes from an AS to a neighbouring AS’, is described in section 10 item (b) of RFC [RFC4364]. In the forwarding plane for Multi-AS Backbones Option B, an ASBR forwards packets based on the MPLS label-value associated with each VPN-IPv4 route.

The remainder of this Annex specifies policy features for addressing operational requirements of the MPLS-ICI without changing the procedures of [RFC4364].

### **A.1 Routing**

BGP is specified in [RFC4364] as the routing protocol used across service providers boundaries, e.g. MPLS-ICI. This section discusses specific additional policies and mechanisms that must be supported across the MPLS-ICI.

#### **A.1.1 Route Target Allocations – Multi-AS Backbones Option B**

An ASBR in compliance with this Technical Specification must provide an SP with the ability to define import/export policies between two MP-BGP peers that apply to the Route Target (RT) community attribute(s) of labeled VPN-IPv4 routes being exchanged over the corresponding MP-BGP session. In addition to supporting the policies defined in the BGP/MPLS VPN MIB [RFC4382], the following policy must be supported:

An SP must be able to export labeled VPN-IPv4 routes that match specified RT values, but with possibly different RT(s) specified as part of the policy. This RT mapping function should be capable of being performed on a one-to-one, one-to-many, many-to-one, or many to many basis. An example of a ‘many to many’ RT mapping policy could be where a single prefix has three RT values assigned. The ASBR will match all three RT values and replace them with three new RT values.

Labeled VPN-IPv4 routes are identified by, amongst other things, the Route Target (RT) extended community attribute. In [RFC4364] Multi-AS Backbones Option ‘B’, labeled VPN-IPv4 routes along with associated RTs are exchanged between service provider domains. Service providers however usually

have difficulty in provisioning RT values that are assigned by external parties on their PE network elements. This difficulty is due to the provider having automated or manual IP VPN service provisioning systems and operational processes that apply to single provider domains only. In order to simplify operation, an ASBR in compliance with this Technical Specification must be capable of mapping an incoming RT value to an outgoing RT value as specified earlier based on import/export policies.

An ASBR may receive a labeled VPN-IPv4 route with associated RT(s) from both internal and external peers and should be capable of re-advertising the route with one or more mapped RT value(s). This mapping function should be supported with labeled VPN-IPv4 routes received from both an MPLS-ICI and an internal peer.

### **A.1.2 Route Distinguisher (RD) Allocations – Multi-AS Backbones Option B**

An ASBR in compliance with this Technical Specification must provide an SP the capability to configure the following policies that apply to the RD part of a labeled VPN-IPv4 route:

- Reject routes whose RD matches specific values.
- Reject routes whose RDs match a specified type, administrator subfield, or assigned number subfield.

In addition, an ASBR should provide an SP the capability to configure an RD rewrite policy as follows:

- Upon exporting a route, replace the RD part of the route with an RD value specified in the policy. The policy should allow replacing the full 8-byte RD value, the administrator subfield and/or the assigned number subfield.

In [RFC4364] Multi-AS Backbones Option B labeled VPN-IPv4 routes are exchanged among ASs via MP-BGP [RFC4364]. A VPN-IPv4 address is composed of an 8-byte Route Distinguisher (RD) pre-pended to a 4-byte IPv4 address. An RD pre-pended to an IPv4 address must uniquely define the VPN-IPv4 address. When labeled VPN-IPv4 routes are exchanged across provider domain boundaries, VPN-IPv4 addresses must not overlap. In order to ensure that RD values that transit the MPLS-ICI are unique and do not collide with those used inside the AS of the receiving ASBR, RD values must in principle contain a registered AS Number (ASN) [RFC4364]. To enforce that, an ASBR in compliance with this Technical Specification that receives VPN-IPv4 routes from another provider-AS must be capable of rejecting routes whose RD or part of their RD matches a provider-specified value. For instance, an SP may want to reject routes that it receives from a peering SP if their administrator subfield value matches its own ASN.

An ASBR must in addition be capable of re-writing the RD value within a VPN-IPv4 address based on SP specified policies applicable to both received and advertised routes. When a public ASN is used as part of the encoding of type 0 or type 2 RD values [RFC4364], RD rewrite should only involve changes to the 2 or 4 byte administrator subfield, leaving the assigned number subfield unchanged. This assignment policy should be carefully coordinated amongst service providers for VPNs that will transit AS borders.

RD rewrite may be invoked to provide for: (1) security and confidentiality, (2) simplified management of the Inter-provider IP VPN, (3) diversified routes through alternative ASBRs and potentially load-balancing when Equal Cost Multipaths (ECMPs) exist, or (4) enabling traffic routing across different ASBRs to the destination based on geographic location. This rewrite function ensures that for security purposes, VPN-IPv4 addresses are unique to each service provider domain and that they do not reveal any numbering scheme used within the provider network. In addition, RD rewrite provides for PEs within an



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SP domain to learn reachability to the same IP VPN destination in the peering provider domain or beyond through different ASBRs that have MP-BGP peering sessions with that peering provider. A PE may select an ASBR as the exit point to a peering provider based on the cost to reach the ASBR, or may load-balance traffic across multiple ASBRs when the paths through these ASBRs have equal costs.

### **A.1.3 Route Target Advertisement – Multi-AS Backbones Option B**

The ASBR should be capable of restricting RT values to only be advertised to and received from interested parties across interconnects. This is particularly relevant where the ASBR is peered with more than one other ASBR. [RFC4684] describes this function across ASes. In order to hide the intra-domain IP VPN topology from an external service provider, RTs advertised across the MPLS-ICI should consist of different values from the RT's provisioned within a service provider's domain.

### **A.1.4 VPN-IPv4 Route Summarization**

An ASBR must be capable of summarizing VPN-IPv4 routes learned from or advertised across the MPLS-ICI boundary on a per VPN basis [RFC4364]. Per VPN route summarization allows existing operational processes for reducing numbers of customer routes being advertised across a single provider's network to apply at the AS border. In Multi-AS Backbones Option A of [RFC4364], route summarization must be configurable at an ASBR on a per BGP session basis in the context of a Virtual Routing and Forwarding (VRF) table. For Multi-AS Backbones Option B of [RFC4364], route summarization must be configurable at an ASBR on a per VPN basis. In the case of Option B, summarization should be practiced with caution as an advertised label across the MPLS-ICI must map to a label associated with the PE that advertised the summarized VPN-IPv4 routes within the advertising ASBR AS.

## **A.2 Label Value Acceptance – Multi-AS Backbones Option B**

For Inter-provider IP VPN based on [RFC4364] Multi-AS Backbones Option B, MP-BGP is used to exchange labels for VPN-IPv4 routes. The LSP segment established via BGP across the MPLS-ICI is a segment of an LSP that extends from PE to PE across ASes. As per paragraph 2 of Section 10 (b) of RFC 4364, "An ASBR should never accept a labeled packet from an EBGP peer unless it has actually distributed the top label to that peer". However, the VPN-IPv4 label may not be the top label on the received labeled packet as the packet may be tunneled by another LSP (e.g., TE-LSP) and the VPN label is the bottom label in the stack. An ASBR must never accept a labeled packet from an EBGP peer unless it has actually distributed all labels exposed at the ASBR to the peer. These include the top arriving label and every label exposed after popping a label. An ASBR should log the number of packets discarded as a result of this processing for diagnostic purposes.

## **A.3 DSCP/Precedence-EXP Mapping**

An ASBR that supports Multi-AS Backbones Option A as specified in [RFC4364] must enable the following:

- Type 1 map: Configuration of a map between incoming IP Diffserv Code Point values (DSCP) and outgoing EXP field values, and the association of that map with an MPLS-ICI or a logical channel over the MPLS-ICI (i.e., an IP VPN attachment circuit [RFC4364]). Upon receiving an IP packet over an MPLS-ICI, an ASBR must be able to classify the IP packet based on DSCP and determine the corresponding EXP value according to the configured map. The ASBR must be able to write the

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outgoing EXP value in the outer MPLS header and possibly inner header (based on configuration) when sending the resulting MPLS labeled packet on one of its core interfaces

- Type 2 map: Configuration of a map between incoming IP Precedence values and outgoing EXP field values, and the association of that map with an MPLS-ICI or a logical channel over the MPLS-ICI (i.e., an IP VPN attachment circuit [RFC4364]). Upon receiving an IP packet over an MPLS-ICI, an ASBR must be able to classify the IP packet based on Precedence bits and determine the corresponding outgoing EXP value according to the configured map. The ASBR must be able to write that EXP value in the outer MPLS header and possibly inner header (based on configuration) when sending the resulting MPLS labeled packet on one of its core interfaces.
- Type 3 map (applies to Ethernet 802.1q interfaces): Configuration of a map between incoming EXP field value to an ASBR and priority bits in the 802.1q VLAN Tag and the association of that map with an MPLS-ICI or a logical channel over that MPLS-ICI (e.g., the VLAN ID that acts as an attachment circuit identifier to a VRF [RFC4364]). Such map may be simply to copy the EXP value to the 3 priority bits in the VLAN tag. An ASBR must be able to apply that map and write the priority bits in the VLAN tag upon encapsulating the IP packet in an 802.1q Ethernet frame prior to transmission to the next ASBR across the MPLS-ICI.
- Type 4 map (applies to Ethernet 802.1q interfaces): Configuration of a map between incoming priority bits carried in an 802.1q attachment circuit VLAN tag and outgoing EXP value and the association of that map with an MPLS-ICI or a logical channel (attachment circuit) over that MPLS-ICI. Upon receiving an IP packet encapsulated in an 802.1q Ethernet frame, the ASBR must be able to apply that map, and accordingly write the EXP value in the MPLS label header(s) it imposes on the IP packet prior to transmission towards its AS-bound core interfaces.

An ASBR, that supports Multi-AS Backbones Option A as specified in [RFC4364] should enable one or more of the following:

- Type 5 map: Configuration of a map between incoming EXP field values and outgoing DSCP values, and the association of that map with an MPLS-ICI or a logical channel over the MPLS-ICI (i.e., an IP VPN attachment circuit [RFC4364]). Upon receiving an MPLS packet over an interface, an ASBR must be able to classify the received MPLS packet based on outer label EXP field value and determine the corresponding outgoing DSCP value according to the configured map. The ASBR must be able to re-write that DSCP value in the IP header when sending the packet over the MPLS-ICI. This map type does not maintain transparency to customer IP DSCP marking. Maintaining transparency is often a customer requirement.
- Type 6 map: Configuration of a map between incoming EXP field values and outgoing precedence values, and the association of that map with an MPLS-ICI. Upon receiving an MPLS packet over an interface, an ASBR must be able to classify the received MPLS packet based on outer label EXP field value and determine the corresponding outgoing precedence value according to the configured map. The ASBR must be able to re-write that precedence value in the IP header when sending the IP packet over the MPLS-ICI. This map type does not maintain transparency to customer IP precedence marking.
- Type 7 map: Configuration of a map between incoming EXP field values plus DSCP value and outgoing DSCP values, and the association of that map with an MPLS-ICI or a logical channel over the MPLS-ICI (i.e., an IP VPN attachment circuit [RFC4364]). Upon receiving an MPLS packet over a core interface (i.e., an AS-internal link), an ASBR must be able to classify the received MPLS

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packet based on outer label EXP field value + DSCP of the encapsulated IP packet and determine the corresponding outgoing DSCP value according to the configured map. The ASBR must be able to re-write that DSCP value in the IP header when sending the IP packet over the MPLS-ICI. This map type does not maintain transparency to customer IP DSCP marking.

- Type 8 map: Configuration of a map between incoming EXP field values + IP Precedence value and outgoing precedence values, and the association of that map with an MPLS-ICI or a logical channel over the MPLS-ICI (i.e., an IP VPN attachment circuit [RFC4364]). Upon receiving an MPLS packet over a core interface (i.e., an AS-internal link), an ASBR must be able to classify the received MPLS packet based on the outer label EXP field value + IP Precedence of the encapsulated IP packet and determine the corresponding outgoing precedence value according to the configured map. The ASBR must be able to re-write that precedence value in the IP header when sending the IP packet over the MPLS-ICI. This map type does not maintain transparency to customer IP precedence marking.
- Type 9 map: Configuration of a map between incoming EXP field values + IP Precedence value and outgoing DSCP values, and the association of that map with an MPLS-ICI or a logical channel over the MPLS-ICI (i.e., an IP VPN attachment circuit [RFC4364]). Upon receiving an MPLS packet over a core interface (i.e., an AS-internal link), an ASBR must be able to classify the received MPLS packet based on outer label EXP field value + IP Precedence of the encapsulated IP packet and determine the corresponding outgoing DSCP value according to the configured map. The ASBR must be able to re-write the DSCP value in the IP header when sending the IP packet over the MPLS-ICI. Specifically, when writing the DSCP value, the 3 bits corresponding to IP precedence must remain untouched. This map type maintains transparency to customer IP precedence marking if the ASBR on the other end implements Type 10 map.
- Type 10 map: Configuration of a map between incoming DSCP value and outgoing precedence value and EXP, and the association of that map with an MPLS-ICI or a logical channel over the MPLS-ICI (i.e., an IP VPN attachment circuit [RFC4364]). Upon receiving an IP packet over an MPLS-ICI, an ASBR must be able to classify the received IP packet based on DSCP and determine the corresponding outgoing IP precedence value according to the configured map. The ASBR must be able to re-write the DSCP value in the IP header with the 3 bits corresponding to IP precedence preserved in the DSCP field as received and the rightmost three bits re-written with 0. This map type, co-jointly with Type 9 map applied at the other end of the MPLS-ICI, maintains transparency to customer IP precedence marking.

#### A.4 Resiliency

For both Multi-AS Backbones Option A and Option B, an ASBR should support the configuration of a multi-hop BGP session with another ASBR between corresponding loopback addresses. When an ASBR interconnects to one or more ASBRs via MPLS-ICIs, failure of one MPLS-ICI should trigger fast traffic re-direction to the surviving interfaces through which the destination is reachable. Implementations should target reroute times on the order of 50-100 milliseconds or better.

When an ASBR peers with two or more ASBRs and an IPv4 route (Multi-AS Backbones Option A) or a VPN-IPv4 route (Multi-AS Backbones Option B) is learned via two or more of its ASBR neighbors, the failure of one ASBR should result in fast traffic re-direction to the surviving ASBR(s). For Multi-AS Backbones Option B, ASBRs must allow the tunneling of VPN-IPv4 labels via MPLS-TE tunnels protected by MPLS fast reroute for link protection.

For Multi-AS Backbones Option B, an ASBR compliant with this Technical Specification must support the rewrite of the RD part of a VPN-IPv4 route received across the MPLS-ICI as described earlier. This allows reachability to an IPv4 destination in a VPN to be learned through two different ASBRs providing for faster convergence.

An ASBR compliant with this Technical Specification must support MP-BGP graceful restart [RFC4724] [RFC4781] for the VPN-IPv4 Address Family Identifier (AFI)/Subsequent Address Family Identifier (SAFI) for Multi-AS Backbones Option B. It must also support BGP graceful restart for the IPv4 AFI in the case of Multi-AS Backbones Option A. Graceful restart enables the data plane to continue forwarding in the presence of control plane failure. It should be noted that graceful restart is effective if the BGP routes in the forwarding plane continue to be reachable through the ASBR that advertised them during graceful restart. Proper operation requires that the ASBR performing graceful restart for eBGP also performs graceful restart for the IGP on the AS-internal interfaces as well as graceful restart for the MPLS protocols (e.g., LDP and/or RSVP-TE) that provide for tunneling the IP VPN packets to the PE connected to the customer site or the far-end ASBR connected to another AS. If the ASBR supports a non-stop control plane (e.g., maintains routing state and signaling state on an active control card and a standby control card, and resumes signaling and routing without interruption after control card switchover), control plane switchover will not be noticed by the rest of the network. In that case, the ASBR need only implement graceful restart in helper mode. Otherwise, the ASBR must implement graceful restart in both helper and restart modes.

#### **A.5 Admission Control Policy: VPN-IPv4 Route Learning Restriction Requirement**

An ASBR must be capable of restricting the number of VPN-IPv4 routes learned on a per VPN basis when supporting [RFC4364] Multi-AS Backbones Option B. These VPN-IPv4 routes could be learned from the service provider's own network or from other service provider domains. Per VPN route capping allows existing operational processes for limiting the number of customer routes advertised across a single provider network to apply at the AS border. For Multi-AS Backbones Option B, an ASBR can identify the routes belonging to a VPN based on RT(s). An ASBR must support policies that cap the number of unique VPN-IPv4 routes learned per VPN on an MP-BGP session basis and peering group basis.

For Multi-AS Backbones Option A, an ASBR must provide the capability of limiting the number of IPv4 routes per VRF on the ASBR. In addition, an ASBR must provide the capability to limit the number of unique routes learned per BGP session and per BGP peering group that applies to multiple BGP sessions potentially with the same peer SP on the ASBR.

The configuration of capabilities discussed in this section must be left to the discretion of the SP operating the ASBR. Specifically, an SP may choose to configure route caps per VPN, per BGP session, or per peering group, or impose no caps.

#### **A.6 MIBs**

The ASBR should support the management information model consistent with the BGP-4 MIB objects defined in [RFC4273],

The ASBR should support a management information model consistent with the full compliance of the following MIB groups as specified in [RFC4382]:

- mplsL3VpnScalarGroup

- mplsL3VpnVrfGroup
- mplsL3VpnIfGroup
- mplsL3VpnPerfGroup
- mplsL3VpnVrfRteGroup
- mplsL3VpnVrfRTGroup
- mplsL3VpnSecGroup
- mplsL3VpnNotificationGroup

An informative description of these objects, their indexing, content, suggested usage and context is contained in Appendix I.

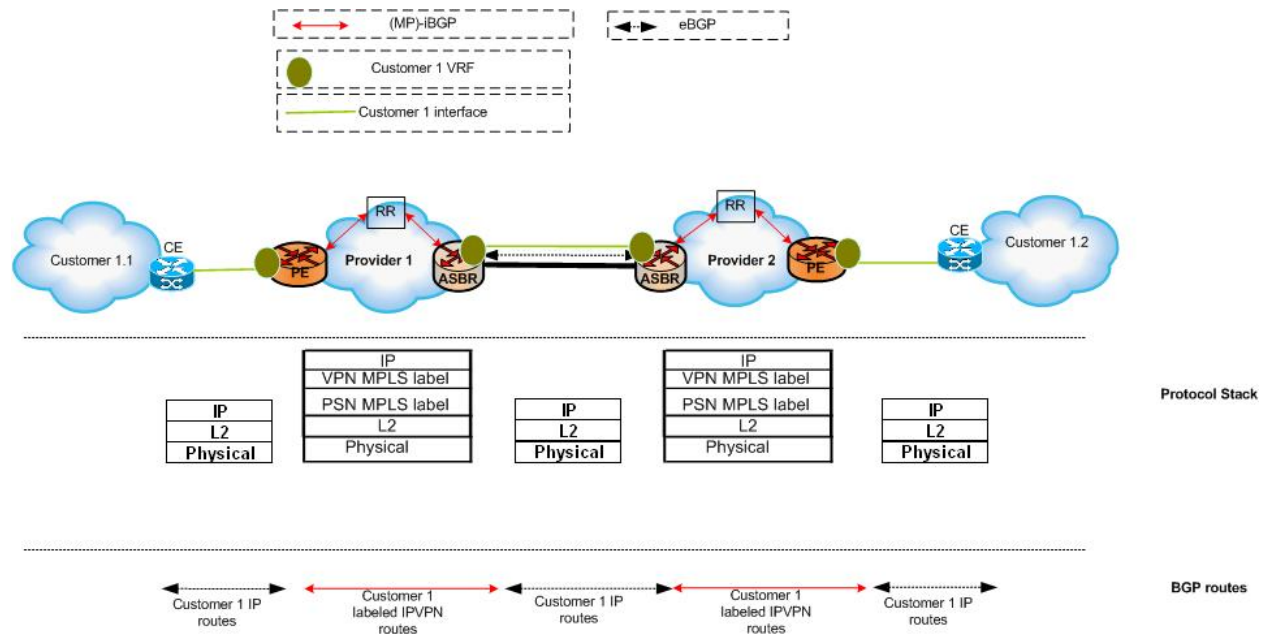
The following BGP/MPLS VPN configuration and management functions are not specified in [RFC4382]. The ASBR should support an information element in the form of enterprise MIB or equivalent to implement the following functions:

The ASBR shall be configurable on a per peer basis with the maximum number of routes that can be received from a peer. The configuration shall allow the user to specify the action taken in the event this threshold is crossed as either sending an alarm message, dropping routes when reaching a number of learned routes that is greater than the configured value, or both of these actions.

#### **A.7 BGP/MPLS IPv4 VPN – Multi-AS Backbones Option A (non-normative)**

(This section does not form an integral part of this Technical Specification)

An Inter-Provider IP VPN service is depicted in Figure 7 using [RFC4364] Multi-AS Backbones Option A. In Option A, described as ‘VRF-to-VRF connections at the AS (Autonomous System) border’, one ASBR in one AS attaches to an ASBR in the other AS via one or more attachment circuits. There is at least one attachment circuit for each IP VPN whose routes are exchanged via these two ASBRs. Each such circuit is an IP sub-interface to an IP VPN VRF configured on the corresponding ASBR. IP VPN Forwarding decisions at the ASBRs are based on IP destination address lookup within the corresponding IP VPN VRF forwarding table. In Multi-AS Backbones Option A, non-labeled IP packets are exchanged between the ASBRs across an MPLS-ICI.

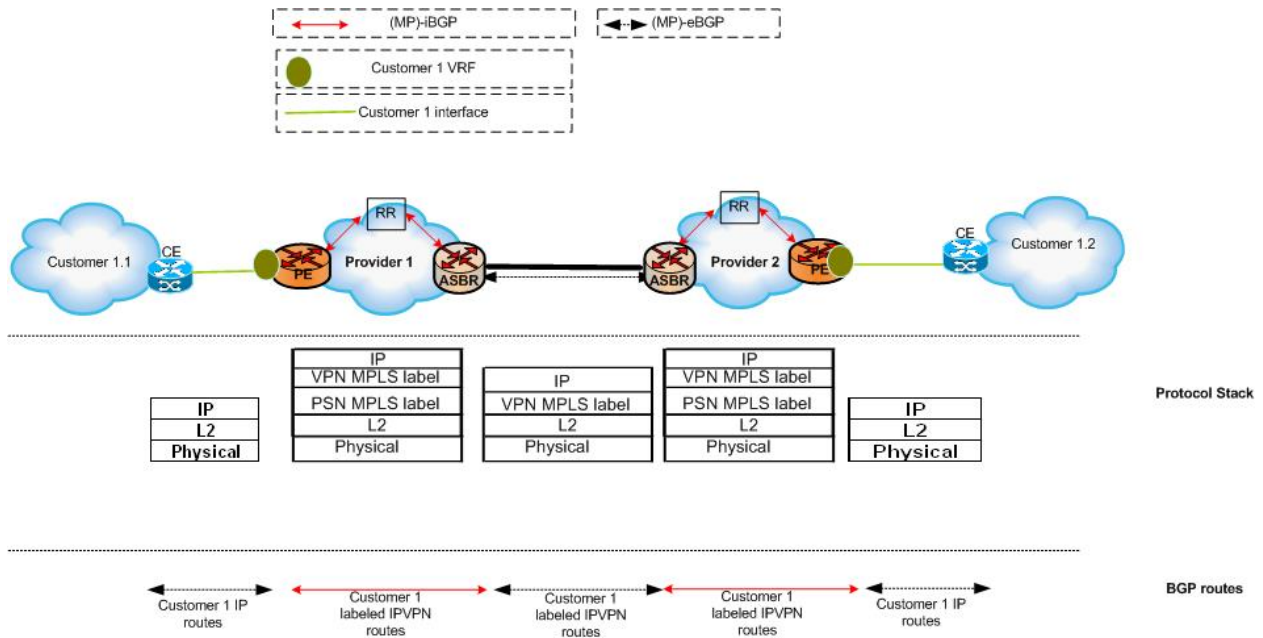


**Figure 7: Inter-provider IP VPN service using RFC 4364 Multi-AS Backbones Option A**

**A.8 BGP/MPLS IPv4 VPN – Multi-AS Backbones Option B (non-normative)**

(This section does not form an integral part of this Technical Specification)

An Inter-Provider IP VPN service is depicted in Figure 8 using [RFC4364] Multi-AS Backbones Option B. In Multi-AS Backbones Option B, described as ‘EBGP redistribution of labelled VPN-IPv4 routes from AS to neighbouring AS’, each ASBR receives labelled VPN-IPv4 routes from within its domain and redistributes these labelled routes to connected ASBRs with which it maintains e-BGP peering sessions supporting labelled VPN-IPv4 routes. When redistributing VPN-IPv4 routes, the ASBR assigns MPLS label values to the routes from its own label space(s). In addition, the ASBR may modify other attributes associated with these routes. In the forwarding plane, an ASBR forwards packets based on the MPLS label-value associated with each VPN-IPv4 route.



**Figure 8: Inter-provider IP VPN service using RFC 4364 Multi-AS Backbones Option B.**

### A.9 Interconnect Performance Measurement (non-normative)

(This section does not form an integral part of this Technical Specification)

Service providers need to measure performance metrics such as round-trip delay, one-way delay, delay variation and packet loss in order to provide for SLA guarantees across an MPLS-ICI. In the case of [RFC4364] Multi-AS Backbones Option A and Option B, these performance metrics can be measured to either designated host addresses within customer VPNs or to an ASBR address for ASBR to ASBR MPLS-ICI data plane performance measurements. Service providers need to have the ability to host measurement probes, as well as carry and report performance measurements across an MPLS-ICI. Guidelines and methodologies for performing QoS measurements and budgeting for impairments across multiple provider domains can be found in [Y.1541] and [MIT\_WP].

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## Annex B Labeled IPv4 Routes

An ASBR that conforms to this Technical Specification must support the distribution of MPLS-labeled IPv4 routes across an MPLS-ICI using MP-BGP. An ASBR that learns labeled IPv4 routes via MP-BGP from another ASBR across an MPLS-ICI may in turn distribute MPLS labels for these routes internal to its AS using an Internal BGP session (iBGP) or other label distribution protocols (e.g., LDP). Doing so allows the establishment of a label switched path to a destination route across an MPLS-ICI. The benefit of distributing a label with an IPv4 route advertised by BGP is that it allows the establishment of LSPs to IPv4 destinations as well as distributes reachability information for IPv4 routes via a single protocol, namely BGP.

### B.1 MP-BGP – Label Distribution

[RFC3107] defines the MP-BGP extensions necessary to distribute one or more labels with an IPv4 route advertised by a BGP update message and to withdraw an advertised labeled IPv4 route. An ASBR that supports this Technical Specification must support [RFC3107].

An ASBR must allow enabling/disabling RFC3107 capability on an MP-BGP session basis. When RFC3107 is enabled on a session, route update processing must be controllable via BGP route policies. At a BGP speaker, a route may be learned with multiple labels and the same BGP next hop. Similarly, the same route may be learned with multiple labels and different BGP next hops. In either case, the receiving ASBR should be able to keep all labels and load balance traffic across equal cost labeled routes. The minimum number of equal cost multipaths that can be supported should be 4.

For security purposes, MP-BGP sessions with labeled IPv4 route advertisements should be run on direct point-to-point links and the BGP next hop for labeled routes must be constrained to be the advertising speaker. All other labeled routes must be rejected at the receiving speaker. In addition, the label stack (there could be one or more labels per stack) must be processed at the BGP next hop. Finally, implementations must support per-interface label spaces. When a labeled packet arrives on an interface and the outer label is outside of that interface space, the packet must be dropped.

### B.2 Routing

MP-BGP with extensions defined in [RFC3107] is the routing protocol used to distribute an IPv4 route and the label associated with that route in addition to other routing information across an MPLS-ICI.

### B.3 Resiliency

The ability to reroute around failure in this case is dependent on BGP convergence. When ECMP entries are available, a system should be capable of removing a failed route from the set of ECMP entries, keeping the other surviving ECMP entries, prior to full BGP route re-selection so that traffic is directed away from the failed path fast enough. A system should also enable tunneling a BGP session over a TE tunnel between the BGP speakers with MPLS FRR enabled on the tunnel for next hop and/or link protection. In the case of MPLS FRR, to maintain security, labels associated with BGP routes should be allocated from a label space restricted to the TE tunnel, treated as an interface. This will allow link/node protection to work while maintaining a level of confidence in the received label authenticity.

An ASBR compliant with this Technical Specification must support BGP graceful restart [RFC4724] and [RFC4781] for the labeled IPv4 route AFI/SAFI in order to continue forwarding for external BGP routes upon control plane switchover. It should be noted that the validity of these routes is dependent on the routes being reachable through the AS of the advertising ASBR. Proper operation requires that the ASBR



performing graceful restart for eBGP also performs graceful restart for the IGP on the AS-internal core interfaces. If the ASBR supports a non-stop control plane (i.e., maintain routing state and signaling state on an active control card and a standby control card, and resume signaling and routing after control card switchover without interruption), control plane switchover will not be noticed by the rest of the network. In that case, the ASBR need only implement graceful restart in helper mode. Otherwise, the ASBR must implement graceful restart in both helper and restart modes.

#### **B.4 Policies and Admission Control**

An ASBR compliant with this Technical Specification must support the definition of MP-BGP import and export policies that control the processing of labeled IPv4 routes. In addition, an ASBR must provide the capability to control the number of labeled routes that could be kept or accepted per BGP session from a BGP speaker or a set of BGP speakers that happen to belong to the same peering AS. In addition, the number of ECMP entries per labeled route must be configurable.

#### **B.5 MIBs**

The ASBR should support the management information model consistent with the BGP-4 MIB objects defined in [RFC4273]. There is no standard MIB for BGP status updates or statistics. Therefore, a router enterprise MIB or equivalent information model is necessary.

The ASBR must be configurable to enable/disable BGP state change notifications.

If BGP is enabled, the ASBR shall generate a status change notification for at least the following events:

- BGP peer transitions from up state to down state.
- BGP peer transitions from down state to up state.
- Number of BGP routes crosses pre-configured threshold.
- Number of BGP route changes over a pre-configured time interval crosses a pre-configured threshold

The ASBR shall be configurable to generate an alarm message if the number of routes received exceeds the threshold value. This shall be configurable per BGP session. The action taken shall be configurable to either continue accepting routes or drop routes once the threshold is exceeded.

## Annex C Pseudowires

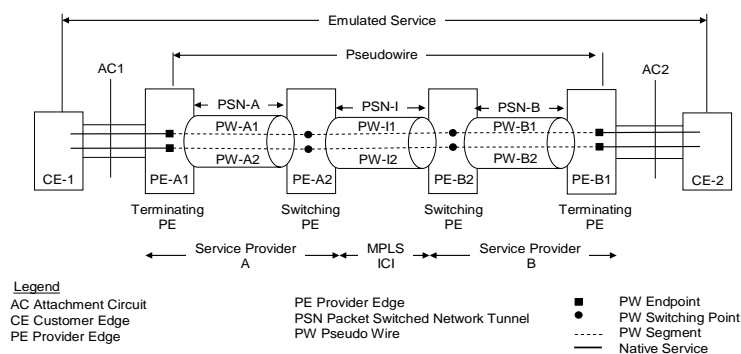
### C.1 Statically-configured and dynamically signaled MPLS Pseudowires

#### C.1.1 Signaling

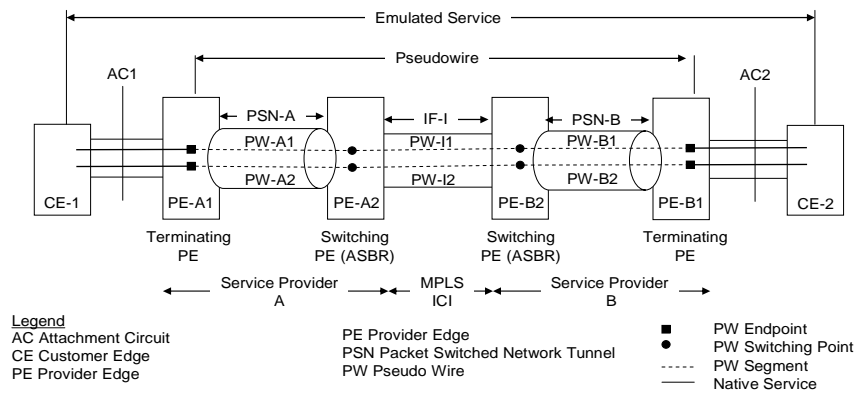
Figure 9 shows the reference architecture in support of the MPLS-ICI MS-PW service, adapted from [MS-PW\_Requirements] to the MPLS-ICI configuration and terminology, where the PW segment (e.g., PW-I1) is carried over the MPLS-ICI over a PSN tunnel (e.g., PSN-I). The PW reference architecture [RFC3985] specifies a PSN tunnel that interconnects the provider edges (PEs) at the terminating ends of the PW segment and tunnels that PW segment. The objective of the PSN tunnel is to hide the MPLS PW label exchanged between the PEs from the PSN, specifically the intermediate nodes on the path between these PEs. When the PEs terminating the PW segment are directly connected, as in the case of the PW segment extending between two ASBRs (PEs) directly connected by an MPLS-ICI, there is no need to have a PSN tunnel that tunnels the PW over the direct link (the MPLS-ICI). As stated in Section 1 of RFC 4447, “packets that are transmitted from one end of the pseudowire to the other are MPLS packets which must be transmitted through an MPLS tunnel. However, if the pseudowire endpoints are immediately adjacent and penultimate hop popping behavior is in use, the MPLS tunnel may not be necessary.”

Figure 10 shows the reference architecture in support of the MPLS-ICI MS-PW service, adapted from [MS-PW\_Requirements] to the MPLS-ICI configuration and terminology, where the PW segment (e.g., PWI1) over the MPLS-ICI is carried over the MPLS-ICI directly without a PSN tunnel.

An ASBR compliant with this Technical Specification must support [RFC4447] for PW signaling (e.g., PW-A1, PW-I1 and PW-B1). For the establishment of multi-segment static switching of PW segments at ASBRs (i.e., Switching PEs PE-A2 and PE-B2 in Figure 9 and Figure 10) that interconnect service provider domains (e.g., Service Provider A and Service Provider B in Figure 9 and Figure 10), the PW segment over the MPLS-ICI is established using PW control LDP signaling [RFC4447].



**Figure 9: Multi-Segment PW Reference Model with a PSN Tunnel at the MPLS-ICI**



**Figure 10: Multi-Segment PW Reference Model without a PSN Tunnel at the MPLS-ICI**

### C.1.2 Connection Admission Control

An ASBR compliant with this Technical Specification must support admission control for PWs based on traffic parameters described in Section 10 for statically configured LSPs (each direction of a PW is an LSP). The traffic parameters and CoS of the PW segment over the MPLS-ICI must be configurable at the segment endpoints (ASBR at both ends of the MPLS-ICI) since they are not carried in signaling messages.

### C.1.3 Routing

An ASBR at the head of a PW segment that spans an MPLS-ICI needs to know how to reach (i.e., over what interface) the ASBR IP address at the other end of the PW segment over the MPLS-ICI. This IP reachability can be learned via BGP or static routing. An implementation compliant with this Technical Specification may support either method.

### C.1.4 Resiliency

An ASBR compliant with this Technical Specification must support

- LDP graceful restart [RFC3478] in helper mode for PW ID FEC 128 [RFC4447].
- LDP graceful restart [RFC3478] in helper mode for PW ID FEC 129 [RFC4447] if the ASBR implements FEC 129 for PW signaling.
- BGP graceful restart [RFC4724] in helper mode if eBGP is used to determine reachability to the ASBR at the other end of the MPLS-ICI.

- 
- Graceful restart for the MPLS tunneling protocol used to tunnel the PW if the LDP PW segment is tunneled over an MPLS PSN tunnel
  - Graceful restart in restart mode for
    - LDP with above mentioned LDP FECs,
    - BGP
    - MPLS tunneling protocol(s).

if the failure of its control plane disrupts the LDP session(s), eBGP sessions or the PSN tunnels (e.g., RSVP-TE tunnels) with its neighbors,

If the ASBR supports a non-stop control plane (e.g., maintain routing state and signaling state on an active control card and a standby control card, and resume signaling and routing after control card switchover without interruption), control plane switchover will not be noticed by the rest of the network. In that case, the ASBR need only implement graceful restart in helper mode.

#### **C.1.4.1 PW Protection Models**

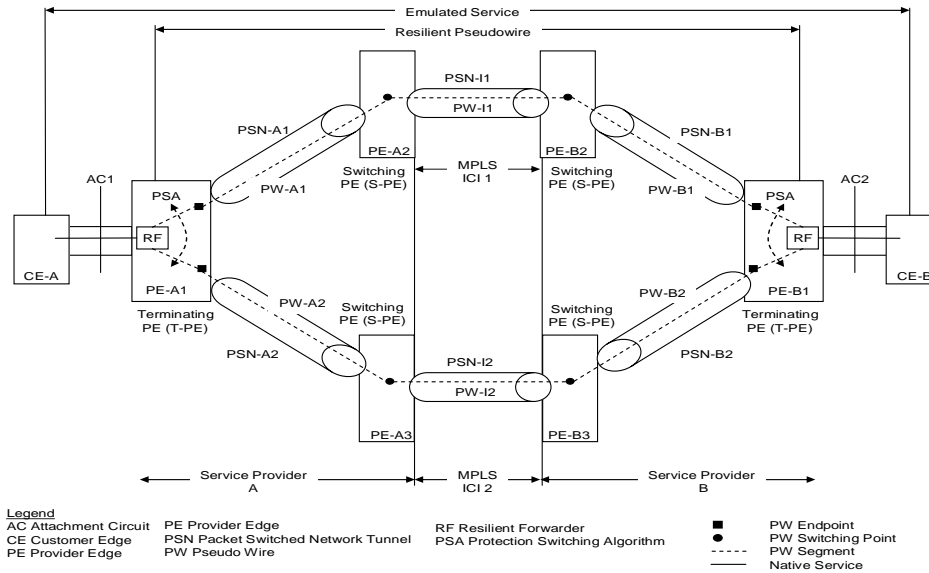
This section covers the following protection models:

- End-End MS-PW Protection. The ASBR role in this case is solely failure detection and notification.
- MPLS-ICI protection for a PW segment over the MPLS-ICI

##### **C.1.4.1.1 Support for End-to-End Protection Model**

###### **C.1.4.1.1.1 ASBR Role**

In the end-to-end protection model depicted in Figure 11, an ASBR (S-PE) with a PW segment set up over an MPLS-CI is responsible for detecting failures on the MPLS-ICI and for subsequently sending failure notification to the ASBR on the other end of the MPLS-ICI and to the PW segment within its own AS towards the T-PE. The ASBR (S-PE) is also responsible for detecting the failure of the PW segment within its own AS, terminated on the ASBR, and for sending failure notification to the other end of the segment and the ASBR on the other end of the MPLS-ICI. The ASBR is also responsible for relaying failure notifications it receives from its own AS or from the ASBR on the other end of the MPLS-ICI toward the T-PE in the direction of the received failure notification. Failure detection and notification procedures performed by the ASBR are defined in section C.2 which covers PW OAM.



**Figure 11: ASBR Fault Detection/Notification at the MPLS-ICI used in end-to-end protection (non-normative)**

#### C.1.4.1.1.2 T-PE Role (non-normative)

(This section does not form an integral part of this Technical Specification)

In the end-to-end protection model, T-PEs should receive the MS-PW failure notification and take action to switchover traffic to an alternative path. However, T-PE behavior is out of the scope of this Technical Specification.

A T-PE may implement PW path protection or PW protection as discussed in this section. Figure 11 depicts an end-to-end protection model whereby a MS-PW has a primary path through one MPLS-ICI and a redundant path through another MPLS-ICI. In both cases, a PW segment is set up between the ASBRs interconnected by these MPLS-ICIs. MPLS-ICI failure detection and failure notification to the T-PEs is the responsibility of the ASBRs, while traffic switchover between primary and redundant paths is the responsibility of the T-PEs.

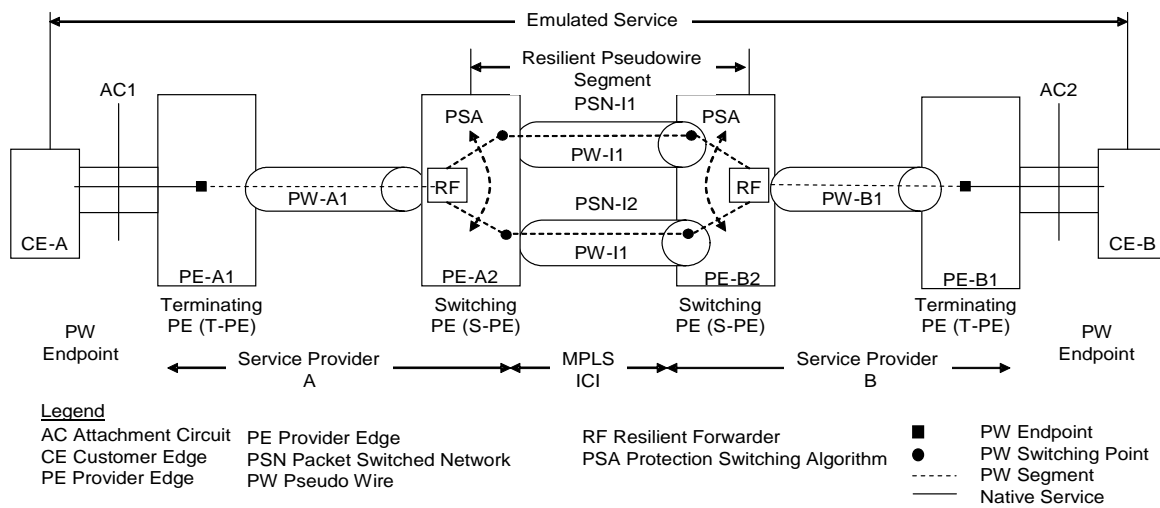
The paths need only be specified in terms of T-PEs and S-PEs. The primary path has at least one PW segment set up between two ASBRs (S-PEs) interconnected by an MPLS-ICI while the redundant path has at least one PW segment set up over a diverse MPLS-ICI. It should be noted that in the context of this Technical Specification, an MS-PW is a concatenation of PWs whereby at least one PW segment is set up over an MPLS-ICI.

The selection between the primary path and redundant path is the responsibility of each T-PE with an endpoint of the MS-PW. Since both directions of the MS-PW must follow the same path, each T-PE must be consistently configured with which path is primary and which path is redundant and must consistently elect to send traffic on one of the two paths from either direction. When there is no failure, T-PEs must select the primary path for sending traffic. Signaling extensions that allow each T-PE to indicate the primary path and indicate switchover provide operational enhancements but are not yet defined, and such definition is out of scope of this Technical Specification.

In order to support switchover from a primary path to a redundant path, an ASBR that detects the failure of a PW segment must be able to send a failure notification towards the T-PEs to trigger switchover from the primary path to the redundant path or vice-versa. Since there is no Type-Length-Value (TLV) defined yet for signaling to indicate which path is being used by a T-PE or which path has a fault, in-band OAM via VCCV can be used to inform T-PEs of the identification of the faulty path.

Alternatively, PW protection may be used. Two different PWs (different FEC128 PWIDs, or (SAII, TAI)) can be used for primary and redundant paths albeit they must be associated with the same Attachment Circuit at the T-PEs. In that case, only one PW can be up at a time. When a T-PE signals to the remote T-PE a PW FEC associated with the same attachment circuit, but different from the one currently used, it should be taken by that T-PE as an indication to switchover. In this latter case, LDP procedures for failure notification can be used. During the transition from one path to the other, T-PEs may have elected different paths due to differential delays in receiving and processing OAM messages. LDP failure notification may use status notification messages [RFC4447] or label withdrawal messages if status message is not supported. This specification recommends that equipment implement the status notification message.

**C.1.4.1.2 MPLS-ICI Protection for a PW segment over the MPLS-ICI**



**Figure 12: MS-PW Segment Protection over Parallel MPLS-ICIs**

An ASBR, in conformance with this specification must support local PW segment protection. In local PW segment protection, there will be a primary path and a redundant path for a MS-PW segment between the same two ASBRs. The path may be a direct MPLS-ICI or a tunnel between the same ASBR pairs

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traversing one or more MPLS-ICIs, as defined in section C.1.1. Figure 12 depicts the case when the MS-PW segment primary and redundant paths are set up over parallel MPLS-ICIs between the same ASBR pair where a PSN tunnel is present. When two ASBRs are interconnected by parallel MPLS-ICIs, and all that is needed is link protection, the redundant path should not have to be configured. That is, an ASBR that detects the failure of the PW segment over an MPLS-ICI must attempt to automatically select an alternative MPLS-ICI that satisfies the PW requirement for sending the PW traffic on. If the redundant path is pre-set up (i.e., the alternative interface is pre-selected), the ASBR must be able to switch traffic to the alternative interface as soon it detects the failure (i.e., it does not have to make an admission control decision). If the PW segment between two ASBRs is set up over a tunnel, tunnel reroute will cause seamless reroute of the carried PW segment. It should be noted that if the ASBRs are not interconnected by parallel MPLS-ICIs or if an ASBR may need to reroute a PW over a multi-hop path to the ASBR at the other end of the PW segment, the two ASBRs must be interconnected by a tunnel over the new path that tunnels the rerouted PW segment between the same ASBR pair. Segment protection between the same pair of ASBRs may be preferred over end-to-end path protection for efficiency reasons and/or the speed in seamless local reroute around a link failure.

For MPLS-ICI PW segment protection, there are four protection models:

1. PW segment without a tunnel – PW protection is provided by a redundant MPLS-ICI (redundant L1/L2 links) as defined in Section 9.1.1.1.1.

The redundant MPLS-ICI path is provided by a single hop L1/L2 link. The applicable protection model is the one-hop model defined in Section 9.1.1.1 and Figure 4. The MPLS-ICI PW segment can be either statically configured per Section 9.1 or signaled via LDP per section C.1.1.

An ASBR compliant with this Technical Specification shall support this MPLS-ICI PW segment protection model.

2. Multi-hop PW segment protection as defined in Section 9.1.1.1.2 and Figure 5.

A multi-segment PW (either stand alone or nested within MPLS tunnels) can provide the redundant path.

An ASBR should support this MPLS-ICI PW segment protection model

3. Multi-hop PW segment protection over a bypass MPLS tunnel (one in each direction) without FRR as defined in Section 9.1.1.1.3 and Figure 6.

The PW over the MPLS-ICI can be set up over a tunnel or the MPLS-ICI directly. If the tunnel exists, it is not protected via MPLS FRR. The PW itself is rerouted over a bypass tunnel connecting the ASBR pair at both ends of the PW upon MPLS-ICI failure detection.

An ASBR should support this MPLS-ICI PW segment protection model

4. PW segment over MPLS PSN TE-tunnel (one in each direction) – TE-tunnel fast reroute [inter-AS-RSVP-TE] as discussed in Annex D.

The TE-tunnel in this section is effectively the “MPLS LSP” that is protected by the bypass tunnel in Section 9.1.1.1.3. The “bypass tunnel” is the facility bypass tunnel defined in [inter-AS-RSVP-TE]. This facility bypass tunnel protects the TE tunnels and as a result protects the PWs carried over it.

An ASBR should support this MPLS-ICI PW segment protection model

### C.1.4.1.3 Revertive and Optimum PW Protection Switching

PW revertive protection switching behavior described in this section applies to PW protection models 1, 2, and 3 as defined in the previous section. An ASBR should enable the configuration of a primary and redundant path for a PW. In addition, an ASBR should allow PW revertive switching from the redundant to the primary path after protection switching.

Revertive PW switching may cause traffic in transit to arrive at the destination out of order. On the other hand, under certain circumstances, the primary path may be more optimum than the redundant path. In that case, an operator may want to revert the PW to the primary path when it becomes active again. The detection of primary path availability at the T-PE is out of scope of this document. However, an ASBR with an MPLS-ICI that protects its own segment of a MS-PW should be able to revert back to a PW segment primary path, if enabled by configuration,

An ASBR must also enable rerouting the PW segment over a more optimum path if one becomes available. A more optimum path can be one that traverses fewer hops if it is riding on a multi-hop tunnel.

### C.1.4.1.4 Switchover

Upon detecting the failure of a PW segment, an ASBR, with pre-configured and established local protection, must be able to switchover the PW or PW segment to the alternative path in 50-100 msec. This specification is concerned with segments set up over an MPLS-ICI and the ASBRs at either end of that MPLS-ICI, be it S-PEs or T-PEs.

## C.2 PW OAM

[MS\_PW\_Requirements] describes requirements for MS-PW OAM mechanisms. These requirements are applicable to MS-PWs set up across an MPLS-ICI with ASBRs as S-PEs. ASBRs interconnected by an MPLS-ICI can also be Terminating PEs (T-PEs) with native attachment circuits. This section covers PW OAM mechanisms and capabilities that are required for operations across an MPLS-ICI. It addresses both PWs and MS-PWs.

The OAM mechanisms described in this section capitalize on mechanisms described in [vccv]. They require some configuration capabilities on ASBRs interconnected by an MPLS-ICI and some procedures to provide for simultaneous end-to-end and segment-based OAM operations for MS-PWs. This Technical Specification is limited in scope to segment-based OAM as it applies to an MPLS-ICI as a segment on a PW path. The objectives of the OAM mechanisms described in this section are to provide for:

- PW segment failure detection and notification to other segments of a MS-PW
- Transparent MS-PW OAM support end-to-end across the network. S-PEs (ASBRs) must be able to pass T-PE to T-PE PW OAM messages transparently
- Support for switchover between 1:1 protected PWs end-to-end.
- Support for switchover between 1:1 protected PW segments across an MPLS-ICI

The MPLS-ICI may represent an un-trusted boundary between providers. In such cases, it may be undesirable for an ASBR in one provider network to respond to certain PW OAM messages originating in



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another provider's network and received on an MPLS-ICI. Therefore, ASBRs when acting as S-PEs for PWs, should support the ability to discard PW OAM messages received on an MPLS-ICI. Procedures described in this section provide for such capability.

## C.2.1 PW OAM Mechanisms

### C.2.1.1 PW Connection Verification: “always on” failure detection and notification

As specified in this Technical Specification, the PW segment established across the MPLS-ICI is statically configured but could be statically established or dynamically established using LDP PWE3 control [RFC4447]. In addition, end-to-end OAM may or may not be configured for the MS-PW.

In cases where end-to-end OAM is not enabled for the MS-PW and the segment across the MPLS-ICI is statically established, VCCV [vccv] becomes the only available mechanism to check the liveness of the MS-PW segment across the MPLS-ICI. The ASBR must relay the status of the PW to other segments by interworking with LDP.

An ASBR compliant with this Technical Specification must support the following capabilities:

- In-band VCCV [vccv] based on the PW associated channel for the VCCV Control Channel (CC) traffic as specified in [vccv]
- BFD with IPv4/UDP header encapsulation for fault detection and status notification as the Continuity Verification type (CV). The CV type is carried in the control word [vccv]. It is either configured on both ends of a PW or negotiated via targeted LDP when targeted LDP is enabled on the MPLS-ICI for PW setup.
- Interworking between BFD status message and LDP status notification message as discussed in the failure detection and notification section C.2.2.
- Asynchronous BFD for continuity check. The following must be configurable for a BFD session at the ASBRs on either end of the PW segment:
  - IP address used as the source IP address when sending BFD messages and as the destination IP address when receiving BFD message
  - The transmission interval for generating the BFD messages must be configurable. In addition, the failure detection time must be configurable as a multiple of the transmission interval from the peer.
  - Min receive interval
  - PHB assigned to labeled packets that contain BFD messages. Operators should ensure continuity check messages receive the most assured and lowest latency forwarding behavior available so that false failure detection does not occur.
- The ASBR may allow the configuration of :
  - Your Discriminator [BFD\_base]
  - My Discriminator [BFD\_base]

- If a BFD session fails, the ASBR that detects the failure must send a failure notification to the remote end of the session via a BFD status message. It must also map that failure to the preceding segment via LDP status notification. The status code used in the status notification message must indicate the failure of the PW segment and is to be assigned by the IETF.
- If the ASBR is a T-PE for a PW, the ASBR should map the failure notification to an OAM mechanism native to the attachment circuit

When LDP and VCCV are used to indicate status of a PW segment, the resolution of the status of a PW segment must follow [vccv] Section 4.1. In cases where the ASBR-ASBR PW segment is statically set up, there is no LDP to indicate failure of a PW across an MPLS-ICI.

An ASBR compliant with this Technical Specification must allow for end-to-end OAM and segment OAM simultaneously. In particular, within the scope of this specification is the PW segment at the MPLS-ICI. An ASBR compliant with this Technical Specification must enable testing the liveness of a PW segment across the MPLS-ICI and failure notification or status up messages (i.e., “PW forwarding” status code from [RFC4446] meaning all faults are cleared) across such a segment. In order to provide for these capabilities, the following must be provided on the ASBR:

- The ability to manually configure a BFD session per PW/PW segment on each ASBR at either end of that segment across the MPLS-ICI
- The ability to set the PW label TTL value for each BFD message generated by an ASBR to 1 and the TTL value in the IP header of the encapsulated BFD message to 255. An ASBR that receives a PW MPLS Packet Data Unit (PDU) with TTL 1 will have TTL expire and intercept the PDU for further processing. The ASBR should perform the following checks in processing the PDU. The more of these checks that are done in the forwarding plane without impacting forwarding engine performance, the more resilient the receiving ASBR will be to potential DoS attacks. These checks may be implemented in any order on the local system. The order of the checks as expressed below represents an example.
  - The ASBR processes the packet if the PW label is a valid label in accordance with section 13 (i.e., the label was assigned by the ASBR to the neighbor). Otherwise, the packet is dropped and a corresponding counter is incremented.
  - The ASBR checks if OAM is enabled for that segment. If it is, it proceeds to the next step. Otherwise, the ASBR discards the PDU and increases a counter for TTL-expired PW PDUs.
  - If the PDU is not dropped, the ASBR checks if there is a control word. If there is, processing proceeds with the next step. Otherwise, the packet is dropped.
  - If the CV type matches the configured type for that PW, processing continues. Otherwise, the packet is dropped.
  - If the CV type is that of BFD with IP encapsulation and fault detection, the TTL value in the IP header is checked. If TTL is less than 255, the packet may be assumed to have been compromised and the packet is dropped. Otherwise, processing of the packet continues.
  - If the PDU contains an authentication header, the message is authenticated. If authentication succeeds, processing of the packet continues. Otherwise, the packet is dropped and a counter is

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incremented for dropped OAM packets that failed authentication. This counter should be maintained per PW.

- If the PDU is not dropped, the ASBR checks the IP address of the source. If the source IP addresses matches that of the peer, processing of the packet proceeds with the next step. Otherwise, the packet is dropped.
- If the PDU is not dropped, the ASBR checks that the destination IP address corresponds to itself and is what was used for that BFD session. If this check passes, processing proceeds. Otherwise, the PDU is dropped.
- If the PDU is not dropped, the ASBR checks that the “Your Discriminator” field in the message matches the locally configured “My Discriminator” field and that “My Discriminator” field in the message matches the locally configured “Your Discriminator”. If this check passes, processing proceeds. Otherwise, the PDU is dropped.

In order to allow for simultaneous operation of end-to-end OAM and segment-based OAM as described in this section, end-to-end OAM messages must have the PW MPLS label TTL value set to 255. This avoids expiry of the PW label TTL at an S-PE.

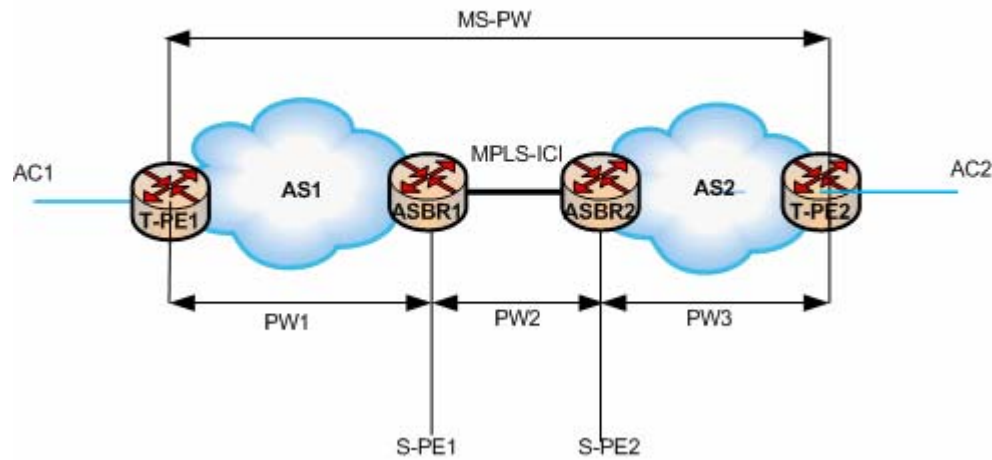
### C.2.1.2 Diagnostics

Diagnostics allows on-demand path verification, checking of control plane against data plane state and for path trace. An ASBR compliant with this Technical Specification must support the following:

- VCCV [vccv] with CV type being LSP-ping and CC being the PW associated control channel.
- LSP-ping in ping mode as specified in [vccv] Section 5.4.1.
- The ability to initiate LSP ping test. An ASBR must be able to test the PW segment extending to the immediate neighbor across the MPLS-ICI by setting TTL to 1 in the PW label header.
- The ability of conducting an LSP-ping test within the provider network. Details of the operations of this test are outside of the scope of this document.

A MS-PW may have multiple switching points within one provider network. A provider may not want to reveal the number of segments traversed by an MS-PW within its domain to another provider. One simple approach to deal with this problem is that an ASBR that receives an LSP-ping echo request across an MPLS-ICI does not reply back with the downstream label map. This behavior must be configurable. However, this behavior will prevent the trace initiator from initiating an LSP-ping test in ping mode to a point beyond the ASBR as it will not be able to include the correct label stack.

## C.2.2 Failure Detection and Notification Procedures



**Figure 13: Reference model for a MS-PW across an MPLS-ICI**

An ASBR compliant with this Technical Specification must support PW failure detection on the MPLS-ICI or an attached PW segment within its own AS. An ASBR, upon failure detection, must send a failure notification to the preceding and the following PW segments or Attachment Circuit (AC) as applicable. In addition, an ASBR must be able to relay failure notification it receives from one PW segment in the forward direction towards the destination T-PE. An ASBR must support PW failure detection and notification for the following connection models in reference to Figure 13:

1. PW1 and PW3 signaled using targeted LDP [RFC4447] and PW2 manually configured
2. PW1, PW2 and PW3 signaled using targeted LDP [RFC4447]
3. PW1 signaled using targeted LDP [RFC4447], PW2 and PW3 manually configured
4. PW1 and PW2 manually configured and PW3 signaled using targeted LDP [RFC4447].

For manually configured PW segments, BFD VCCV [vccv] must be used as the failure detection and notification mechanism. When connecting a signaled segment to a manually configured segment, an ASBR must be able to translate failure notification received via LDP status notification to a BFD diagnostic code and vice versa.

When all segments of a PW are signaled, an ASBR must propagate the status failure notification it receives from one segment to another in the direction of the received notification using LDP. An ASBR, in this case, must use the PWID (in case of FEC 128) corresponding to the segment on the ASBR when it sends the status notification for that segment.

An ASBR must also support failure notification suppression so that if the state of the PW is down because of a failure or a previous failure notification, a newly received failure notification message does not trigger an additional failure notification. If the earlier failure clears, the more recent failure notification is used to trigger a new failure notification message using BFD or LDP as applicable, until all failure conditions are cleared. When all failure conditions are cleared for an attached PW segment, an ASBR must signal status up for that segment.

### C.3 MIBs

The ASBR should support the management information model consistent with the following MIB groups as specified in [RFC3815] in read-write mode:

- mplsLdpGeneralGroup
- mplsLdpNotificationsGroup

An informative description of these objects, their indexing, content, suggested usage and context is contained in Appendix I.

The ASBR should support the management information model consistent with the following MIB groups as specified in [PWMIB] in read-write mode:

- pwBasicGroup
- pwNotificationsGroup
- pwIDGroup
- pwGeneralizedFECGroup
- pwPwStatusGroup
- pwAttachmentGroup
- pwPerformanceIntervalGroup and pwHCPPerformanceIntervalGroup
- pwSignalingGroup
- pwNotificationControlGroup

An informative description of these objects, their indexing, content, suggested usage and context is contained in Appendix I.

The router PW fault detection times should be configurable via an enterprise MIB for use in mapping to emulated service status indication protocols (e.g., TDM OAM, FR LMI, ATM OAM, Ethernet OAM).

The ASBR should support the management information model consistent with the following MIB groups as specified in [PWMPLSMIB] in read-write mode:

- pwMplsGroup
- pwMplsOutboundMainGroup and pwOutboundTeGroup
- pwMplsInboundGroup
- pwMplsMappingGroup

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An informative description of these objects, their indexing, content, suggested usage and context is contained in Appendix I.

Since PW segment protection is not currently specified in [PWMPLSMIB], an information model for an enterprise MIB should be supported to implement the configuration of protection models specified in section C.1.4.1.2:

- Configuration of specific MPLS-ICIs (with or without tunnel) as primary and redundant in model 1
- Configuration of PW segment as primary and a MS-PW as redundant in model 2.
- Configuration of an MPLS-ICI (with optional tunnel identified in the LSR MPLS XC MIB) as primary and configuration of a redundant MPLS tunnel in model 3
- Configuration of an MPLS-ICI (with TE tunnel identified in the MPLS TE MIB) as primary and a redundant MPLS tunnel in model 4

The ASBR should support an information model where the redundant "path" in the protection model can be statically configured consistently with the protection model.

The ASBR should support an information model where the redundant "path" in the protection model can be either statically configured or dynamically determined (i.e., via RSVP-TE or LDP signaling of a tunnel), consistent with the protection model.

The ASBR should support standard information models for reporting statistics and alarm notifications for the objects in the PW protection models (i.e., MPLS-ICI, PW, non-TE tunnel, TE tunnel).

The ASBR should support an information model for configuration of whether the T-PE local protection switching action should revert to the primary PW segment if it is determined to be available via either LDP status notification and/or VCCV continuity checking as described in section C.2.1.2. Note that this must be consistently configured at each ASBR at the end of a protected PW segment to ensure correct operation.

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## **Annex D Traffic Engineering (TE)-Tunnels**

### **D.1 Statically-configured and signaled TE-tunnels**

#### **D.1.1 Signaling**

An ASBR conformant to this Technical Specification must provide the capability of configuring a TE-tunnel E-LSP to another ASBR over an MPLS-ICI. The LSP itself can be the only LSP segment (i.e., the segment starts at one ASBR and terminates at another ASBR at the other end of the MPLS-ICI), or it can be a segment of an end-to-end LSP initiated in one provider domain and terminated in another, transiting the MPLS-ICI.

An ASBR must have the capability of establishing a configured TE-tunnel to a connected ASBR over an MPLS-ICI via RSVP-TE signaling. An ASBR must provide the capability of stitching, by administrative action, that LSP to another LSP in each provider domain to form an LSP that crosses the domain boundary.

An ASBR must be able to limit the reception of certain label values to a group of one or more interfaces connected to a designated set of peer ASBRs. If an MPLS packet arrives on one of those interfaces with a label that is not associated with that set of peers, the packet must be dropped.

An ASBR must also provide the capability of configuring a traffic profile for the TE-tunnel LSP extended over the MPLS-ICI. The traffic profile should include the bandwidth per class type (CT) (indicating a CoS) supported over the LSP. Section 11.1 addresses the requirements for traffic management in more details.

When dynamically establishing LSP segments across an MPLS-ICI, it is preferred that any intermediate segment establishment be triggered by RSVP-TE signaling, starting at the head-end of the end-to-end LSP.

RSVP-TE signaling must support signaling of multiple CTs for each TE-tunnel E-LSP based on the extensions defined in [MPLS-CNI].

#### **D.1.2 MIBs**

The ASBR should support the management information model consistent with the following read-write MIB groups as specified in [RFC3812]:

- mplsTunnelGroup
- mplsTunnelScalarGroup
- mplsTunnelManualGroup
- mplsTeNotificationGroup

An informative description of the objects in these groups, their indexing, content, suggested usage and context is contained in Appendix I.

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## D.2 Dynamically Established TE-Tunnels

### D.2.1 Signaling

This section states the required and recommended capability for establishing inter-carrier traffic engineering (TE) tunnels that carry traffic for one or more service classes. TE tunnels are LSPs with constraints (e.g., bandwidth, preemption priority, etc.). In particular, this section defines the capabilities required for dynamic setup of TE tunnels across service provider domain boundaries.

An ASBR must be capable of supporting the setup of a TE E-LSP from one router in one service provider domain to another router in the other service provider domain. Currently, the simplest way to achieve this is by extending existing intra-area and intra-domain mechanisms for RSVP-TE, including capability of signaling Diffserv Traffic Engineering (DS-TE) LSPs [RFC4124].

ASBRs compliant with this Technical Specification and this section must be capable of establishing dynamic inter-domain TE-LSPs using interdomain RSVP-TE extensions as defined in [inter-AS-RSVP-TE]. Specifically, an ASBR must support three types of inter-AS TE tunnels: (1) contiguous, (2) stitched and (3) nested [inter-AS-RSVP-TE]. In addition, ASBRs compliant with this Technical Specification must support TE-path computation as defined in [interdomain-PDPC]. In this section, stitching is referred to as dynamic stitching. Dynamic stitching, in contrast to static stitching in the control plane, refers to stitching individually signaled RSVP-TE sessions together in the data plane and control plane.

### D.2.2 Routing

This Technical Specification requires that an ASBR support per domain path computation procedures (PDPC) to compute the path segment to the first hop in a received or configured Explicit Route Object (ERO) or to the TE tunnel destination. At an ASBR at the ingress to an MPLS-ICI, the ASBR must be able to select the neighbor ASBR and associated link that satisfies the TE constraints and provide for reachability to the first hop in the ERO and the TE tunnel tail-end. At an ASBR on the egress from an MPLS-ICI, the ASBR must be able to compute a path (generally a partial path) that satisfies the TE constraints and can reach the first hop in the ERO and TE tunnel tail-end.

The only routing information required in this Technical Specification to be exchanged over the MPLS-ICI dynamically is that of reachability relayed via BGP update messages. That is, this specification does not require any change to normal BGP operations between two ASBRs.

### D.2.3 Resiliency

An ASBR in compliance with this Technical Specification must support MPLS-TE Fast Reroute (FRR), providing for fast reroute around MPLS-ICI link failure and ASBR node failure. This Technical Specification requires support for a bypass tunnel, that can tunnel one or more protected TE LSPs, for MPLS-TE FRR. An ASBR in compliance with this Technical Specification should support one-to-one detours for MPLS-TE FRR. It should be noted that providing for node protection requires revealing the node ID next to the ASBR in the forwarding direction as well as the label assigned by that node. This may be considered by some carriers as revealing confidential information, and as a consequence an operator may not allow this behavior and thus FRR may not be allowed.

In order to provide for continued forwarding when the ASBR control plane fails, an ASBR compliant with this Technical Specification must support RSVP-TE graceful restart [RFC3471] as well as graceful restart for the routing protocols on which RSVP-TE path computation depends (i.e., eBGP for IPv4 routes [RFC4724] on the MPLS-ICI and IGP-(TE) on intra-AS core interfaces). If the ASBR supports a non-



stop control plane (e.g., , maintains routing state and signaling state on an active control card and a standby control card, and resumes signaling and routing after control card switchover without interruption), control plane switchover will not be noticed by the rest of the network. In that case, the ASBR need only implement graceful restart in helper mode. Otherwise, the ASBR must implement graceful restart in both helper and restart modes.

#### D.2.4 MIBs

The ASBR should support the management information model consistent with [RFC2206] for provisioning an RSVP session per destination, including at least the following: Sender IP address (rsvpSenderAddr), Destination IP address (rsvpSessionDestAddr), and RSVP message refresh interval (rsvpSenderInterval).

The ASBR should also support the following additional RSVP session parameters using an Enterprise MIB and CLI: Enable/disable RSVP Refresh Reduction [RFC2961], RSVP session status, LSP and Tunnel IDs provisioned by a particular session.

The ASBR should support injection of a TE-LSP into the local Routing Information Base (RIB) using a configured metric via an enterprise MIB or equivalent configuration.

The ASBR should support an enterprise MIB or equivalent that supports configuration of the MPLS PSN tunnel fault detection times for each of the OAM protocols configured to detect faults (e.g., L1, L2).

The ASBR should support the management information model consistent with the following read-write MIB groups as specified in [RFC3812]:

- mplsTunnelGroup
- mplsTunnelScalarGroup
- mplsTunnelSignaledGroup
- mplsTeNotificationGroup

An informative description of the objects in these groups, their indexing, content, suggested usage and context is contained in Appendix I.

### D.3 Admission Control

An ASBR that supports admission control in compliance with this Technical Specification must provide the following admission control capabilities:

- When administratively stitching TE-LSPs at an ASBR, one LSP segment may be signaled up to the pre-configured ASBR endpoint where it is administratively stitched to another LSP segment set up over an MPLS-ICI. If the arriving traffic profile in the RSVP-TE path message differs from that configured for the TE-LSP on the ASBR MPLS-ICI, the ASBR must be capable of rejecting the LSP setup (i.e., the Path message) by sending a PathErr message to the source of the path message.
- When dynamically stitching TE-LSPs, an ASBR must be able to apply admission control on the arriving path message. If the traffic profile of the arriving LSP per Class Type (CT) exceeds that of the LSP, over the MPLS-ICI, with which stitching is required, the LSP setup is rejected via a PathErr

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message. If the arriving path message has bandwidth requirements that are below those of the configured LSP, the LSP is considered to pass the bandwidth admission control. Once admitted, no other LSP can be stitched to the same LSP segment set up over the MPLS-ICI.

- When establishing a contiguous LSP, the egress ASBR at a domain boundary must have the capability of applying admission control on the arriving path message. If the requested bandwidth per CT exceeds the available bandwidth over each of the MPLS-ICI(s) leading to the destination and possibly constrained by the next hop in a received ERO, the ASBR must reject the LSP setup.
- When establishing a contiguous LSP, the egress ASBR at a provider domain boundary must have the capability of applying admission control on the arriving path message. If the requested bandwidth per CT exceeds the available bandwidth allocated for the ASBR AS across the MPLS-ICI(s) leading to the next AS in the ERO or dynamically selected to reach the destination, the ASBR must reject the LSP setup.
- When establishing a nested TE-LSP, the ASBR at a domain boundary must have the capability of applying admission control on the arriving path message. If the requested bandwidth per CT exceeds the available CT bandwidth on the LSP that will be used to tunnel the LSP being requested, the LSP is rejected. Otherwise, the LSP is admitted and the tunnel LSP bandwidth is appropriately adjusted for the corresponding CT(s). ASBRs should provide options for allowing the expansion of the tunnel LSP bandwidth per CT to accommodate the LSP being tunneled. If the expansion is successful, the LSP will be successfully admitted into the tunnel.
- In addition to making admission based on bandwidth per CT, an ASBR must be able to make admission control decisions based on other types of policies. An ASBR that receives TE-LSP setup requests must enable the configuration of policies that apply to admission control. Specifically, an ASBR must be able to make admission control decisions based on:
  1. requested setup priority
  2. requested pre-emption priority
  3. request for route recording
  4. destination address of the LSP
  5. source address of the LSP
  6. neighbor ASBR
  7. affinity of the LSP

An ASBR should be able to admit or reject the setup of a TE-LSP, or modify any of the attributes in the setup (i.e., Path) message based on configuration. If an attribute is modified based on a policy, the source should be notified of the modification.

Policies may be described by a set of one or more match criteria. Policies may also specify a list of one or more associated actions for each match criterion.

The match criteria are applied to RSVP-TE Path messages. If an LSP meets the match criteria, then the associated actions are applied to that LSP. Match may be based, for example, on any combination of:

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1. Interface
  2. Neighbor (for point to point links this is implied by the interface, but for multipoint links there may be multiple neighbors per interface)
  3. Direction (LSP is incoming or outgoing on the associated interface)
  4. Source Address (or prefix)
  5. Destination Address (or prefix)
  6. Setup and/or Holding Priority
  7. Affinity

#### **D.4 Protection of Confidential Information**

An ASBR that supports RSVP-TE signaling over an MPLS-ICI must allow operators to configure policies that control the handling of EROs and Route Record Objects (RROs) for confidentiality reasons.

In general, an ASBR that supports this Technical Specification may have some interfaces which support MPLS-ICIs with “untrusted” neighbors, and some interfaces which support MPLS-ICIs with “trusted” neighbors (e.g., service providers that are known to operate their networks in a relatively secure fashion). The amount of trust may vary based on whether the interface is authenticated, or other factors. Similarly, it may be necessary for security and availability reasons to restrict LSPs from neighbors which fail to implement RSVP Refresh Reduction [RFC2961]. In many cases policies may vary between neighbors based on factors such as these.

The following description outlines one way that policies may be realized. This is not intended to require any particular implementation, and any implementation which provides equivalent functionality is permitted.

An ASBR that supports RSVP-TE signaling over an MPLS-ICI should allow operators to set policies which limit which EROs will be accepted in incoming RSVP-TE Path messages. Policies should be configurable on a per-interface basis. For example, policies may allow operators to prohibit EROs, or to restrict which addresses may be included in EROs (e.g., such as prohibiting addresses of core routers internal to the service provider, while allowing addresses of PE routers or external addresses).

An ASBR that supports RSVP-TE signaling over an MPLS-ICI should allow operators to set policies restricting the transmission of RROs. Policies should be configurable on a per-interface basis. For example, policies may prohibit transmission of RROs (implying that RROs would be stripped out prior to transmission of Path messages), or may prohibit the transmission of the addresses of internal core routers (so that the internal addresses would be stripped out of the RRO prior to transmission across an MPLS-ICI).

An ASBR that supports RSVP-TE signaling over an MPLS-ICI should allow operators to set policies restricting the transmission of PathErr Messages across the MPLS-ICI. Policies should be configurable on a per-interface basis. For example, PathErr messages received from inside the service provider which are being propagated across the MPLS-ICI may have the error code and/or sub-code changed to avoid exposing internal topology information. Thus, for example the PathErr message may be propagated across the MPLS-ICI with error code=2 (“Policy Control failure”), see [RFC3209] and [inter-AS- RSVP-TE].

Similarly, PathErr messages originated by nodes within the service provider may have the source addresses changed (e.g., to the address of the ASBR supporting the MPLS-ICI) in order to preserve the confidentiality of the nodes internal to the SP.

## **D.5 Class Type Mappings**

When there is a mismatch in the implementation of CoS indicated by Class type (CT) and or TE classes (CT plus priority) between two providers that want to establish an MPLS-ICI, mappings must be administratively defined at each provider's boundary at the time the interconnection is configured. Using these mappings, translation must be performed at the ingress ASBR while signaling the LSP across a provider domain boundary over the MPLS-ICI. ASBRs in compliance with this Technical Specification must be able to perform this translation. Specifically, ASBRs must enable the configuration of an incoming map that defines the mapping of incoming CT/TE-class and PHBs from one domain to outgoing CT/TE-class and PHBs that apply in the next domain.

## **Annex E    Voice Over IP**

The Voice over IP application uses the TE-tunnel transport services discussed in Annex D and does not add any additional requirements on the MPLS-ICI.

Service providers need to have the ability to host measurement probes, as well as carry and report performance measurements across an MPLS-ICI in support of inter-carrier VoIP peering. Probes and performance measurements are beyond the scope of the MPLS-ICI. Guidelines and methodologies for performing QoS measurements and budgeting for impairments across multiple provider domains can be found in [Y.1541] and [MIT\_WP].

## Appendix I Informative Description of MIBs

This Appendix contains informative descriptions of the objects contained in MIB groups specified in normative requirements in the body of this specification. It includes a summary of their naming, indexing, content, suggested usage and context. For further detail and information, the implementer should consult the cited references.

### I.1 MPLS LSR MIB [RFC3813]

The following MPLS MIB tables apply to the configuration of static MPLS LSPs or for use in queries of MPLS cross connects that occur as a result of dynamic MPLS signaling: These tables can be used to provision and manage point-to-point, point-to-multipoint, or multipoint-point MPLS LSP segments.

The ASBR uses the `mplsInterfaceGroupTable` in the `mplsInterfaceGroup` as specified in [RFC3813] as a sequence of `MplsInterfaceEntry` objects indexed by `{mplsInterfaceIndex}` for the purpose of defining and querying the status of MPLS-enabled interfaces. This table includes information regarding the label range that can be received or transmitted on an interface level basis, total bandwidth of the interface, and whether per platform or per interface label space is used.

The ASBR uses the `mplsInterfacePerfGroupTable` in the `mplsPerfGroup` as specified in [RFC3813] as a sequence of `MplsInterfacePerfEntry` objects as an extension to the `mplsInterfaceEntry` table for the purpose of tracking statistics regarding label usage, lookup failures, and fragmented packets on MPLS-enabled interfaces.

The ASBR uses the `mplsTunnelPerfTable` in the `mplsInterfaceGroup` as an augment to the `MPLSTunnelTable` as specified in the MPLS TE MIB of RFC3812 for measuring tunnel performance. This table includes 32- and 64-bit counters of the number of packets forwarded by the tunnel, 32- and 64-bit counters of the bytes forwarded by the tunnel, and a 32-bit counter of the packets dropped because of errors or other reasons.

The ASBR uses the `mplsInSegmentTable` in the `mplsInSegmentGroup` as specified in [RFC3813] as a sequence of `mplsInSegmentEntry` objects indexed by `{mplsInSegmentIndex}`. This table includes the IP address of the previous hop, incoming interface, incoming label, whether label popping should occur, the creator of the cross connect (e.g., manual, LDP, or RSVP-TE), an index to the traffic parameters, and an index to a cross-connect table.

The ASBR uses the `mplsInSegmentPerfTable` in the `mplsInSegmentGroup` as specified in [RFC3813] as a sequence of `mplsInSegmentPerfEntry` objects as an augment to `mplsInSegmentEntry` for monitoring the statistics of an incoming MPLS LSP segment. This table includes 32-bit counters for packets, octets, errors and discards in the receive direction of this segment and a 64-bit counter of octets received.

The ASBR uses the `mplsOutSegmentTable` in the `mplsOutSegmentGroup` as specified in [RFC3813] as a sequence of `mplsOutSegmentEntry` objects indexed by `{mplsOutSegmentIndex}`. This table includes identification of the IP address of the next hop, the outgoing interface, outgoing label, whether a label push(es) should occur and a pointer to a label stack table entry, the creator of the cross connect (e.g., manual, LDP, or RSVP-TE), an index to the traffic parameters, and an index to a cross-connect table.

The ASBR uses the `mplsOutSegmentPerfTable` in the `mplsInterfaceGroup` as specified in [RFC3813] as a sequence of `mplsOutSegmentPerfEntry` objects as an augment to `mplsOutSegmentEntry` for monitoring

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the statistics of an outgoing MPLS LSP segment. This table includes 32-bit counters for packets, octets, errors and discards related to packets destined for this segment and a 64-bit counter of octets sent.

The ASBR uses the `mplsXCTable` in the `mplsXCGroup` as a sequence of `mplsXCEntry` objects indexed by `{mplsXCIndex, mplsXCInSegmentIndex, mplsXCOutSegmentIndex}` as specified in [RFC3813]. This table includes the indices to the incoming segment(s) and outgoing segment(s), an LSP ID, the creator of the cross connect (e.g., manual, LDP, or RSVP-TE), the administrative and operational status of the cross-connect.

The ASBR uses the `mplsLabelStackTable` in the `mplsLabelStackGroup` as a sequence of `mplsLabelStackEntry` objects indexed by `{mplsLabelStackIndex, mplsLabelStackLabelIndex}` as specified in [RFC3813]. This table includes a list of label values to be pushed beneath the top label.

The ASBR uses the `mplsLsrNotificationsGroup` as specified in [RFC3813] that indicates whether an MPLS cross-connect (XC) is up or down.

## I.2 Differentiated Services MIB [RFC3289]

The ASBR uses the `diffServDataPathTable` of the `diffServMIBDataPathGroup` as specified in [RFC3289] as a sequence of `diffServDataPathEntry` objects indexed by `{ifIndex, diffServDataPathIfDirection}` for the purposes of specifying the interface index, direction (i.e., ingress or egress) and a pointer to the first Diffserv component of the components configured on this router and interface.

The ASBR uses the `diffServClfrElementTable` of the `diffServMIBClfrGroup` as specified in [RFC3289] as a sequence of `DiffServClfrEntry` objects for the purposes of identifying the classifier employed, the parameters of the classifier, the precedence order in which it is applied, and a pointer to the next Diffserv component if that classifier applied to the packet is matched.

[RFC3289] specifies L3/L4 classification parameters in the `diffServMultiFieldClfrTable` of the `diffServMIBClfrGroup` as a sequence of `diffServMultiFieldClfrEntry` objects indexed by a row pointer that contains Source/Destination IPv4/IPv6 address prefixes, DSCP, and source/destination transport layer port numbers. The ASBR may support this MIB to allow configuration of Access Control Lists (ACLs) that result in only admitting certain types of packets to the ASBR. This MIB can be used to support the security requirements specified in section 13.

The ASBR uses the `diffServMeterTable` of the `diffServMIBMeterGroup` as specified in [RFC3289] as a sequence of `diffServMeterEntry` objects indexed by a row pointer, which in turn contain row pointers to the parameters of the meter, as well as row pointers to the next Diffserv component to be invoked if the meter succeeds or fails the conformance test as specified in section 3.3.1 of [RFC3289].

The ASBR uses the `diffServTBParamTable` of the `diffServMIBTBParamGroup` as specified in [RFC3289] as a sequence of `diffServTBParamEntry` objects indexed by a row pointer that defines the Token Bucket (TB) parameters for a simple Token bucket, SrTCM both color-blind and color-aware, and TrTCM both color-blind and color-aware. Additional values may be specified in other MIBs.

The ASBR uses the `diffServActionTable` of the `diffServMIBActionGroup` as specified in [RFC3289] as a sequence of `diffServActionEntry` objects indexed by a row pointer, which in turn contain row pointers to the parameters of the action, as well as a row pointer to the next Diffserv component to be invoked as specified in section 3.4.1 of [RFC3289].

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[RFC3289] defines tables for specific Diffserv-related actions for counting, DSCP (re) marking, algorithmic dropping and random dropping as described below.

The ASBR uses the diffServCountActTable of the diffServMIBCounterGroup as specified in [RFC3289] as a sequence of DiffServCountActEntry objects indexed by row pointer that contains 64-bit counters of packets and octets.

The ASBR uses the diffServDscpMarkActTable of the diffServMIBDscpMarkActGroup as specified in [RFC3289] as a sequence of diffServDscpMarkActEntry objects indexed by a row pointer which replaces the DSCP value in the packet with the value stored in the table.

The ASBR uses the diffServAlgDropTable of the diffServMIBAlgDropGroup as specified in [RFC3289] as a sequence of diffServAlgDropEntry objects indexed by row pointer. This table contains the type (e.g., tailDrop, headDrop, randomDrop, alwaysDrop), a row pointer to the next Diffserv component, identification of the queue, the threshold applied to that queue, a pointer to a table of parameters specific to the drop algorithm, and 64-bit counters for the number of packets and octets dropped for the algorithmic and random cases.

The ASBR uses the diffServRandomDropTable of the diffServMIBRandomDropGroup as specified in [RFC3289] as a sequence of diffServRandomDropEntry objects indexed by row pointer. This table contains parameters specific to the random drop algorithm, i.e., Weighted Random Early Detection (WRED) algorithm for the minimum threshold packets and bytes, maximum threshold packets and bytes, maximum drop probability, and the weighting for past queue history and the sampling interval for the WRED algorithm.

The ASBR uses the diffServQTable of the diffServMIBSchedulerGroup as specified in [RFC3289] as a sequence of DiffServQEntry objects indexed by a row pointer. This table contains a pointer to an entry in the diffServSchedulerTable, and row pointers to the minimum and maximum rate tables that contain the values that the scheduler should use.

The ASBR uses the diffServSchedulerTable of the diffServMIBSchedulerGroup as specified in [RFC3289] as a sequence of diffServSchedulerEntry objects indexed by a row pointer. This table contains a pointer to the next Diffserv component, which if it is another scheduler is a way to specify hierarchical scheduling. This table contains the scheduler method (e.g., priority, Weighted Round Robin (WRR), Weighted Fair Queuing (WFQ)), and pointers into the MinRate and MaxRate tables.

The ASBR uses the diffServMinRateTable of the diffServMIBSchedulerGroup as specified in [RFC3289] as a sequence of diffServMinRateEntry objects indexed by a row pointer. This table contains parameters relevant to the scheduler method: a priority value (larger numeric value has higher priority), minimum absolute rate, and a minimum relative rate as a fraction of the interface speed.

The ASBR uses the diffServMaxRateTable of the diffServMIBSchedulerGroup as specified in [RFC3289] as a sequence of diffServMaxRateEntry objects indexed by row pointer and diffServMaxRateLevel. This table contains a maximum absolute rate and a maximum relative rate for a non-work-conserving scheduler (i.e., shaper) as specified in section 3.5.5 of [RFC3289].

### **I.3 Bi-directional Forwarding Detection (BFD) MIB [BFD MIB]**

The ASBR uses the bfdSessTable of the bfdSessionGroup as specified in [BFD MIB] as a sequence of BfdSessEntry objects indexed by {bfdSessIndex} for the purposes of configuring and managing BFD sessions. This table includes the application ID (e.g., MPLS VPN), local and remote discriminator



values, User Datagram Protocol (UDP) port number, session state (admin down, down, initializing, up), a Boolean indication regarding receipt of packets from the remote, operational mode (asynchronous or demand with or without echo), Boolean indications of the local preference for demand and echo modes, Boolean indication of operation through a control plane disruption, session address type (IPv4 or IPv6) and address, the desired minimum transmit interval, minimum receive and echo intervals, the detection time multiplier, and a Boolean indication of local desire to use Authentication.

The ASBR uses the `bfdSessPerfTable` of the `bfdSessionPerfGroup` as specified in [BFD MIB] as a sequence of `BfdSessPerfEntry` objects as an augment to `BfdSessEntry` for the purposes of collecting statistics on each BFD session. This table includes 32-bit counters for the number of BFD messages received and sent, the last time and diagnostic code when communication was lost with the neighbor, the number of times this session has entered the up state since the last ASBR reboot, the last discontinuity time, and 64-bit high-capacity counters for the number of BFD messages received and sent.

The ASBR uses the `bfdSessMapTable` of the `bfdSessionGroup` as specified in [BFD MIB] as a sequence of `bfdSessMapEntry` objects indexed by {`bfdSessApplicationId`, `bfdSessDiscriminator`, `bfdSessAddrType`, `bfdSessAddr`} which contains the value of `bfdSessMapBfdIndex`, as the index into `bfdSessTable` to identify a BFD session between a pair of nodes.

The ASBR uses the `bfdSessNotificationsEnable` of the `bfdNotificationGroup` Boolean scalar to enable/disable the BFD session up and down notifications. A single BFD session up and down notification is issued for a contiguous range of entries in `bfdSessTable` to minimize the emission of a large number of notifications.

The ASBR uses BFD notifications when a new detection time or detection multiplier is negotiated.

The ASBR uses BFD notifications when a change occurs in the BFD session operating mode.

#### **I.4 BGP MPLS L3VPN MIB [RFC4382]**

The [RFC4382] MIB is used in the context of multi-AS Backbones option A and B as detailed in Annex A of this document.

The ASBR uses the following scalar objects from the `mplsL3VpnScalarGroup` of [RFC4382]: number of configured VRFs, number of active VRFs, number of connected VPN interfaces, enabling/disabling all notifications, maximum total number of routers across all VRFs, the minimum interval for issuing the maximum route threshold notification, and a threshold for the number of illegally received labels above which a notification is issued.

The ASBR uses the `mplsL3VpnIfConfTable` of the `mplsL3VpnIfGroup` as a sequence of `MplsL3VpnIfConfEntry` objects as specified in [RFC4382] and indexed by {`mplsL3VpnVrfName`, `mplsL3VpnIfConfIndex`} for the purposes of configuring the VPN attributes of individual interfaces. This includes, IP VPN classification (enterprise, carrier's carrier, or inter-AS), and the route distribution protocol used (e.g., BGP, Routing Information Protocol (RIP), Open Shortest Path First (OSPF)).

The ASBR uses the `mplsL3VpnVrfConfTable` of the `mplsL3VpnVrfGroup` as specified in [RFC4382] and indexed by {`mplsL3VpnVrfName`} for the purpose of configuring and managing a Virtual Routing and Forwarding (VRF) instance. This table includes, the human-readable name (`mplsL3VpnVrfName`), the VPN ID from [RFC4265] which can indicate a specific VPN or all VPNs, operational and administrative status, number of associated interfaces, number of active interfaces, mid- and high-level water marks for the number of routes, and the maximum number of routes allowed.

The ASBR uses the `mplsL3VpnVrfPerfTable` of the `mplsL3VpnPerfGroup` as specified in [RFC4382], which augments `mplsL3VpnVrfTable` to provide performance counters for each VRF. This table includes the number of routes added and the number of routes removed since the last discontinuity, the current number of routes, the number of routes dropped due to exceeding the maximum threshold, and the time of the last discontinuity.

The ASBR uses the `mplsL3VpnVrfRouteTable` of the `mplsL3VpnVrfRteGroup` as specified in [RFC4382] and indexed by {`mplsL3VpnVrfName`, `mplsL3VpnVrfRteInetCidrDestType`, `mplsL3VpnVrfRteInetCidrDest`, `mplsL3VpnVrfRteInetCidrPfxLen`, `mplsL3VpnVrfRteInetCidrPolicy`, `mplsL3VpnVrfRteInetCidrNHopType`, `mplsL3VpnVrfRteInetCidrNextHop`} for the purpose of configuring and managing individual route entries in the VRF. This table includes the destination IP address prefix and length, the next hop IP address prefix and length, the next hop type (e.g., local, remote, black hole, reject), mechanism by which the route was learned, AS number of the next hop, primary and alternate metrics, and an index into `mplsXCTable` [RFC3813] to manually configure the label stack and MPLS cross-connect to be used for this route.

The ASBR uses the `MplsVpnVrfRTTable` of the `mplsL3VpnVrfRTGroup` as specified in [RFC4382] and indexed by {`mplsL3VpnVrfName`, `mplsL3VpnVrfRTIndex`, `mplsL3VpnVrfRTType`} for the purpose of configuring and managing individual Route Target (RT) entries in the specified VRF. This table includes the actual RT and its type (import ,export or both) .

The ASBR uses the following notifications of the `mplsL3VpnNotificationGroup` as specified in [RFC4382]: one interface associated with the VRF is up, one interface associated with the VRF is down, and indication when the mid- or high-level number of threshold is crossed, the number of routes in this VRF has fallen below the high-level threshold, the number of illegal VRF labels exceeds the threshold.

## **1.5 MPLS LDP MIB [RFC3815]**

The following apply to statically configured and signaled via LDP LSPs.

The ASBR uses the `mplsInSegmentLdpLspTable` of the `mplsLdpLspGroup` as specified in [RFC3815] as indexed by {`mplsLdpEntityLdpId`, `mplsLdpEntityIndex`, `mplsLdpPeerLdpId`, `mplsInSegmentLdpLspIndex`} for the purposes of associated an LDP signaled LSP with an MPLS cross connect (XC) as defined in the MPLS LSR MIB of [RFC3813]. This table includes the value of `mplsInSegmentIndex` in `mplsInSegmentTable` of [RFC3813], the label type (generic, FR, ATM), and LSP type (terminating, originating, or cross-connecting).

The ASBR uses the `mplsOutSegmentLdpLspTable` of the `mplsLdpLspGroup` as specified in [RFC3815] as indexed by {`mplsLdpEntityLdpId`, `mplsLdpEntityIndex`, `mplsLdpPeerLdpId`, `mplsOutSegmentLdpLspIndex`} for the purposes of associating an LDP signaled LSP with an MPLS cross connect (XC) as defined in the MPLS LSR MIB of [RFC3813]. This table includes the value of `mplsOutSegmentIndex` in `mplsOutSegmentTable` of [RFC3813], the label type (generic, FR, ATM), and LSP type (terminating, originating, or cross-connecting).

The following requirements apply to targeted LDP-established LSPs.

The ASBR uses the SNMPv3 MIB `mplsLdpEntityGenericLRTable` of the `mplsLdpGeneralGroup` as indexed by {`mplsLdpEntityLdpId` `mplsLdpEntityIndex`, `mplsLdpEntityGenericLRMin`, `mplsLdpEntityGenericLRMax`} from [RFC3815] for assigning a label range to an interface or a platform for the indexed LDP entity.

The ASBR uses the `mplsLdpEntityTable` of the `mplsLdpGeneralGroup` as indexed by `{mplsLdpEntityLdpId, mplsLdpEntityIndex }` of the MPLS LDP MIB [RFC3815] as indexed by `(mplsLdpEntityLdpId, mplsLdpEntityIndex)` to configure/set-up potential LDP sessions on a specific LSR/LER/PE.

The ASBR populates the `mplsLdpPeerTable` of the `mplsLdpGeneralGroup` as a sequence of `mplsLdpPeerEntry` objects indexed by `(mplsLdpEntityLdpId, mplsLdpEntityIndex, mplsLdpPeerLdpId)` of the MPLS LDP MIB [RFC3815] as determined by LDP through initialization or discovery with information about LDP Peers known to LDP Entities through LDP signaling.

The ASBR populates the `mplsLdpSessionTable` of the `mplsLdpGeneralGroup` as a sequence of `mplsLdpSessionEntry` objects indexed by `{mplsLdpEntityLdpId, mplsLdpEntityIndex, mplsLdpPeerLdpId, mplsLdpSessionPeerAddrIndex }` of the MPLS LDP MIB [RFC3815] with information learned about a peer via LDP, such as; state, role, protocol version, keep-alive timer and discontinuity time.

The ASBR populates the `mplsLdpEntityStatsTable` of the `mplsLdpGeneralGroup` as an augment to `mplsLdpEntityTable` of the MPLS LDP MIB [RFC3815] to keep statistical information about the LDP Entities on the ASBR, such as; session-level errors, LDP message content errors, and timer expirations.

The ASBR populates the `mplsLdpHelloAdjacencyTable` of the `mplsLdpGeneralGroup` as a sequence of `mplsLdpHelloAdjacencyEntry` objects indexed by `{mplsLdpEntityLdpId, mplsLdpEntityIndex, mplsLdpPeerLdpId, mplsLdpHelloAdjacencyIndex }` of the MPLS LDP MIB [RFC3815] with statistical information learned about a peer, such as the hold timer and adjacency type.

The ASBR uses LDP Notifications of the `mplsLdpNotificationsGroup` as defined in the MPLS LDP MIB [RFC3815]. The intent of this requirement is to autonomous notification of the management plane of LDP neighbor adjacency changes.

The ASBR uses the `mplsFecTable` of the `mplsLdpGeneralGroup` as specified in [RFC3815] as a sequence of `mplsFecEntry` as indexed by `{mplsFecIndex }` for the purposes of configuring Forwarding Equivalence Class (FEC) information. This table contains the FEC type (prefix or host), the address type (IPv4, IPv6), the address (or prefix), and the prefix length.

## **I.6 Pseudowire MIB [PWMIB]**

The ASBR creates an entry in `pwTable` in the `pwBasicGroup` as specified in [PWMIB] for all locally configured PW types (e.g., Ethernet, TDM, FR, ATM, etc.) to hold generic parameters related to PW creation and monitoring. The ASBR creates corresponding entries in the [PWMPLSMIB] and the specific emulated service MIB as specified by `pwType` object.

The ASBR supports the `pwTable` in the `pwBasicGroup` as specified in [PWMIB] as a sequence of `pwEntry` objects indexed by `{pwIndex }` for the purpose of generic configuration and status monitoring specific pseudowire. This table includes information to specify the PW type, PSN tunnel type, set up/holding priority, address type, local stitching information, attachment IDs, configuration of control word option, MTU size, group ID, inbound and outbound PW label, local and remote operational status and other management information.

The ASBR supports the `pwIndexMappingTable` in the `pwPwIdGroup` as specified in [PWMIB] as a sequence of `pwIndexMappingEntry` objects indexed by `{pwIndexMappingPwType,`

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pwIndexMappingPWID, pwIndexMappingPeerAddrType, pwIndexMappingPeerAddr} for the purpose of performing a reverse mapping only if the basic PW FEC ID is used.

The ASBR supports the pwPeerMappingTable in the pwMappingTablesGroup as specified in [PWMIB] as a sequence of pwPeerMappingEntry objects indexed by { pwPeerMappingPeerAddrType, pwPeerMappingPeerAddr, pwPeerMappingPwType, pwPeerMappingPwID} for the purpose of querying PWs based upon peer, type and pwID.

The ASBR uses configuration control of notifications as specified in the pwNotificationsGroup of [PWMIB] controlling whether up/down and pw deleted notifications are sent, and the maximum overall rate of notifications that the ASBR can send.

The ASBR uses the pwPerfCurrentTable in the pwPerformanceIntervalGroup and pwHCPerformanceIntervalGroup groups as a sequence of pwPerfCurrentEntry objects indexed by {pwIndex} as specified in [PWMIB] for the purpose of reporting per-PW performance information for the current interval. These tables include 32-and 64-bit counters of the number of packets received from the PSN, number of bytes received from the PSN, number of packets sent to the PSN, and number of bytes sent to the PSN.

The ASBR uses the pwPerfIntervalTable in the pwPerformance1DayIntervalGroup as indexed by {pwIndex, pwPerfIntervalNumber} as specified in [PWMIB] for the purpose of per-PW performance information for historical intervals. This table contains up to 96 entries (e.g., 15 minute intervals for a 24-hour period) of the statistics contained in pwPerfCurrentTable, the duration of the interval and an indication of data validity.

The ASBR uses the pwPerfTotalTable in the pwPerformanceGeneralGroup indexed by {pwIndex} as specified in [PWMIB] for the purpose of per-PW performance information for the time since PW establishment or the latest management application reset. This table contains the statistics in pwPerfCurrentTable for the total time since the most recent time that any row counter suffered a discontinuity.

## **I.7 Pseudowire to MPLS PSN MIB [PWMPLSMIB]**

The ASBR uses the pwMPLSTable of the pwMPLSGroup as specified in [PWMPLSMIB] as a sequence of pwMPLSEntry objects indexed by {pwIndex} for the purpose of associating an MPLS PSN tunnel with a specific pseudowire. This MIB table includes information to configure the PSN tunnel type (TE or non-TE), the EXP bits used in the PW label, the PW label TTL value, and the peer LDP ID.

The ASBR uses the pwMPLSOutboundTable of the pwOutboundMainGroup and pwOutboundTeGroup as specified in [PWMPLSMIB] as a sequence of pwMPLSOutboundEntry objects which augments an pwMPLSEntry for the purpose of statically mapping a specific PW to a PSN tunnel in the outbound direction.

The ASBR uses the pwMPLSInboundTable of the pwMplsInboundGroup as specified in [PWMPLSMIB] as a sequence of pwMPLSInboundEntry objects indexed by {pwIndex} for the purpose of indicating the inbound PSN tunnel when LDP signaling is used. The single value, pwMPLSInboundXcIndex indicates the XC index representing this PW in the inbound direction.

The ASBR uses the pwMPLSNonTeMappingTable of the pwMplsMappingGroup as specified in [PWMPLSMIB] as a sequence of pwMPLSNonTeMappingEntry objects indexed by {pwMPLSNonTeMappingDirection, pwMPLSNonTeMappingXcIndex,

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pwMPLSNonTeMappingIfIndex, pwMPLSNonTeMappingVcIndex } for indicating the PSN tunnel under a variety of non-TE conditions. An application can use this table to quickly retrieve the PW carried over specific non-TE MPLS outer tunnel or physical interface for use in test access.

## **I.8 MPLS Traffic Engineering (TE) MIB [RFC3812]**

The ASBR uses the mplsTunnelTable from [RFC3812] of the mplsTunnelGroup as a sequence of mplsTunnelEntry objects indexed by {mplsTunnelIndex, mplsTunnelInstance, mplsTunnelIngressLSRId, mplsTunnelEgressLSRId} for establishing MPLS-TE tunnels and tracking their state. This table identifies the Ingress and Egress LSR IDs, whether the tunnel corresponds to an interface or the node, set up/holding priority, session attributes (e.g., FRR, pinning), FRR local protection mode (facility, detour), up time, operational status, alarm state and other attributes.

The ASBR uses the mplsTunnelResourceTable of the mplsTunnelGroup from [RFC3812] which is a sequence of mplsTunnelResourceEntry objects indexed by {mplsTunnelResourceIndex} for setting up the tunnel resources. This table contains the mean bit rate and maximum burst size for the LSP.

The ASBR uses the mplsTunnelHopTable of the mplsTunnelGroup from [RFC3812] as a sequence of MplsTunnelHopEntry objects indexed by {mplsTunnelHopListIndex, mplsTunnelHopPathOptionIndex, mplsTunnelHopIndex} for specifying strict or loose source routed MPLS tunnel hops.

The ASBR uses the mplsTunnelARHopTable of the mplsTunnelGroup from [RFC3812] as a sequence of MplsTunnelARHopEntry objects indexed by {mplsTunnelARHopListIndex, mplsTunnelARHopIndex} to indicate the individual hops for an MPLS tunnel defined in mplsTunnelTable if it is reported by the MPLS signaling protocol (i.e., the record route option has been configured for a particular tunnel).

The ASBR uses the mplsTunnelCHopTable of the mplsTunnelGroup from [RFC3812] as a sequence of MplsTunnelCHopEntry objects indexed by {mplsTunnelCHopListIndex, mplsTunnelCHopIndex} to indicate the hops for an MPLS tunnel defined in mplsTunnelTable as computed by a constraint-based routing protocol, based on the mplsTunnelHopTable for the outgoing direction of the tunnel.

The ASBR uses the following notifications of the mplsTeNotificationGroup from [RFC3812]: Tunnel Up, Tunnel down, Tunnel rerouted, Tunnel re-optimized. The Tunnel Up/Down traps shall be capable of being enabled/disabled.

The ASBR uses the following scalar objects of the mplsTunnelScalarGroup from [RFC3812]: number of tunnels configured, label distribution protocols configured maximum hops that can be specified for a tunnel, and the maximum rate at which tunnel notifications can be sent.

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## Appendix II      MPLS-ICI Forwarding Behavior and EXP bit Mapping Configuration Example

### II.1 Overview

Forwarding behaviors are defined by several standards bodies. The IETF has defined a set of Differentiated Services (Diffserv) Per Hop Behaviors (PHBs) that are used in the definition of Per Domain Behaviors (PDBs) to meet service objectives [RFC3086]. Similarly, [Y.1541] has defined a set of six recommended QoS classes that define objectives that are mapped to PHBs as summarized in section II.4. The examples in this Appendix assume that the ASBR defines a mapping for MPLS EXP bits to local traffic management forwarding behaviors as a local implementation decision.

In the examples of this Appendix, service providers connecting across an MPLS-ICI agree to a set of forwarding behaviors that they will support. These need not be the same, may be a different number, and may not be defined by the same standards body, if any. A set of these forwarding behaviors for each service provider would be identified by EXP bits for each MPLS-ICI logical interface. As defined in [RFC3270], there are only three EXP bits and hence at maximum 8 PHBs assignable per E-LSP. Associating the EXP mapping with an incoming label map [RFC3270] would effectively extend the number of bits available. An EXP<->PHB mapping function as specified in section 11.1.1 is required to map the forwarding behaviors supported in one network to those supported in the network on the other side of the ICI. An EXP-bit remarking function is needed if these mappings are not identical on the ingress and egress interfaces.

Service providers may agree to EXP-bit and forwarding behavior mappings as described in the examples. These mappings should not be interpreted as a hard requirement, but rather represent examples of the generic mapping and remarking functions specified in section 11.1.1 and section 11.5.

## II.2 Example of Bi-Lateral Agreement Forwarding Behavior and EXP-bit Mappings

The following tables show one example of the bi-lateral agreement forwarding behavior and EXP-bit mapping and remarking for two Service Providers (SPs), SP-1 using the currently defined ITU-T Y.1541 QoS Classes (i.e., 0 through 5) and a descriptive phrase taken from the guidance for these QoS classes from section 5.3.5 of [Y.1541] and SP-2 using IETF Diffserv PHBs. Note that per RFC 3270, EXP marking determines the PHB for E-LSPs. In the interest of brevity, drop precedence was not included in this example. For an example of how drop precedence could be encoded in the EXP bits for E-LSPs, see [AGGRDIFFSRV]. This appendix does not set any requirements on how these forwarding behaviors would be specified or signaled.

In addition to the forwarding behavior and EXP bit mapping and remarking, the service providers may also specify policing, scheduling and/or shaping across the MPLS-ICI as specified in sections 11.1.2 and 11.1.3 using the Diffserv MIB information model described in section 11.5.

**Table 3: Example Usage of EXP↔PHB Mappings at SP-1's ASBR**

SP-1 EXP↔PHB Mapping Incoming from SP-1 Domain		Generic PHB "Name"	SP-2 EXP↔PHB Mapping Outgoing to MPLS-ICI		
EXP	Y.1541 QoS Class, Guidance "PHB"		Queue	Diffserv PHB	EXP
0	0, Real-time, jitter sensitive	Alpha	Q2.1	EF	5
2	2, Signaling	Alpha	Q2.1	EF	5
3	3, Transaction Data, interactive	Gamma	Q2.2	AF41	4
4	4, Low Loss Only	Delta	Q2.3	AF31	3
5	5, Default IP	Epsilon	Q2.4	BE	0

SP-2 EXP↔PHB Mapping Incoming from MPLS-ICI		Generic PHB "Name"	SP-1 EXP↔PHB Mapping Outgoing to SP-1 Domain		
EXP	Diffserv PHB		Queue	Y.1541 QoS Class, Guidance "PHB"	EXP
5	EF	Alpha	Q1.1	0, Real-time, jitter sensitive	0
		Beta	Q1.2	2, Signaling	2
4	AF41	Gamma	Q1.3	3, Transaction Data, interactive	3
3	AF31	Delta	Q1.4	4, Low Loss Only	4
0	BE	Epsilon	Q1.5	5, Default IP	5

**Table 4: Example Usage of EXP↔PHB Mappings at SP-2's ASBR**

SP-2 EXP↔PHB Mapping Incoming from SP-2 Domain		Generic PHB "Name"	SP-1 EXP↔PHB Mapping Outgoing to MPLS-ICI		
EXP	Diffserv PHB		Queue	Y.1541 QoS Class, Guidance "PHB"	EXP
5	EF	EXP5	Q2.1	0, Real-time, jitter sensitive	0
4	AF41	EXP4	Q2.2	3, Transaction Data, interactive	3
3	AF31	EXP3	Q2.3	4, Low Loss Only	4
2	AF21	EXP2	Q2.3	4, Low Loss Only	4
0	BE	EXP0	Q2.4	5, Default IP	5

SP-1 EXP↔PHB Mapping Incoming at MPLS-ICI		Generic PHB "Name"	SP-2 EXP↔PHB Mapping Outgoing to SP-2 Domain		
EXP	Y.1541 QoS Class, Guidance "PHB"		Queue	Diffserv PHB	EXP
0	0, Real-time, jitter sensitive/ 2, Signaling	EXP5	Q2.1	EF	5
3	3, Transaction Data, interactive	EXP4	Q2.2	AF41	4
4	4, Low Loss Only	EXP3	Q2.3	AF31	3
		EXP2	Q2.3	AF21	2
5	5, Default IP	EXP0	Q2.4	BE	0

In this example, the service providers SP-1 and SP-2 agreed to an equivalence mapping of the ITU-T QoS classes and IETF PHBs, which is indicated in the column titled PHB Name in the middle of the pairs of mappings in Table 3 and Table 4. The generic PHB names in SP-1's ASBR refer to some internal tag used within the ASBR. EXP bit remarking would occur on the egress interface in this ASBR. The generic PHB names in SP-2's ASBR refer to EXP bit markings. In SP-2's ASBR, the EXP bits would be remarked on the ingress interface before being sent to the egress interface. All that needs to be agreed is an equivalence mapping to the generic PHB name that meets the QoS objectives agreed to between the providers. For each ASBR there is a table for each direction, the first in the direction incoming from the SP domain outgoing to the MPLS-ICI, and the second in the direction incoming from the MPLS-ICI and outgoing to the SP domain. Note that not all forwarding behaviors supported by one provider may have a corresponding forwarding behavior in another provider network. Each table shows the EXP to PHB mapping used in the incoming direction in the leftmost two columns, the generic PHB Name in the third column and the corresponding queue used, the PHB, and the EXP bit marking for the outgoing direction in the rightmost three columns.

In this example, SP-1 has five forwarding behaviors based on [Y.1541] while SP-2 also has five forwarding behaviors based on Diffserv [RFC2475]. In this example, SP-2 uses five queues, while SP-1 uses only four to illustrate the possibility that many PHBs may map to a single queue, for example, in the first mapping in Table 4. Note that neither SP followed each of these standards exactly in this example,



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which does occur in the current state of the industry. The mapping between Y.1541 QoS Classes and IETF PHBs also has more entries than that contained in Appendix VI of Y.1541 described in section II.4 below. In most cases in this example, the mapping was one-one. Additionally, the cases of many-one and a null mapping input are also illustrated in this example. The first mapping in Table 3 shows an example of a many-one mapping from SP-1 to SP-2 where two forwarding behaviors from SP-1 are mapped to the single forwarding behavior in SP-2's ASBR, specifically; SP-1 has distinct signaling and real-time QoS classes, while SP-2 only has an Expedited Forwarding (EF) PHB. A similar many-one case exists in SP-2's ASBR in the first mapping in Table 4 for the mapping of IETF PHBs for AF3 and AF2 to only ITU QoS Class 4 in SP-1. The second mapping in Table 3 shows an example of null mapping since SP-2 has only the EF PHB and no support for ITU QoS class 2. A similar case exists in the second mapping in Table 4 since SP-1 has no corresponding PHB to AF21 in SP-2.

### II.3 Another Example of Bi-Lateral Forwarding Behavior Mapping

The following tables show another example of the bi-lateral agreement forwarding behavior and EXP-bit mapping and remarking for two Service Providers (SPs). Table 5 shows the mappings employed by SP-A, and Table 6 shows the mappings employed by SP-B. SP-A uses the currently defined (i.e., 0 through 5) as well as provisional (i.e., 6 and 7) ITU-T Y.1541 QoS Classes. The descriptive phrase taken from the guidance of section 5.3.5 of Y.1541 is defined for QoS classes 0-5, but not for 6 and 7. SP-A has also used section 3.2 of RFC 4594 [RFC4594] to define the "network control" generic PHB. In this example, SP-B uses IETF Diffserv PHBs, with the application guidance defined in the paragraph number from [RFC4594] cited to the right of the IETF defined Diffserv PHB acronym in the SP-specific PHB columns of the table.

The format and meaning of these tables is the same as described in the previous example. Also as described above, SP-A and SP-B have agreed on an equivalence to some of the generic PHB names (and hence likely similar queuing behaviors as well), as indicated in the third column of each of the tables. These generic PHB terms are an extension of those described in [AGGRDIFFSRV].

**Table 5: Example Usage of EXP↔PHB Mappings at SP-A's ASBR**

SP-A EXP↔PHB Mapping Incoming from SP-A Domain		Generic PHB "Name"	SP-B EXP↔PHB Mapping Outgoing to MPLS-ICI		
EXP	Y.1541 QoS Class, Guidance "PHB"		Queue	RFC 4594 PHB	EXP
6	2, Signaling, RFC 4594 3.2	Network Control	QB.1	EF, 4.1	5
5	6	Real-Time	QB.1	EF, 4.1	5
3	7	Near Real-Time	QB.2	AF41, 4.5	4
2	3, Transaction Data, interactive	Non Real Time	QB.3	AF31, 4.7	3
1	4, Low Loss Only	Best Effort	QB.4	BE, 4.9	0

SP-B EXP↔PHB Mapping Incoming from MPLS-ICI		Generic PHB "Name"	SP-A EXP↔PHB Mapping Outgoing to SP-A Domain		
EXP	RFC 4594 PHB		Queue	Y.1541 QoS Class, Guidance "PHB"	EXP
		Network Control	QA.1	2, Signaling, RFC 4594 3.2	6
5	EF, 4.1	Real-Time	QA.2	6	5
4	AF41, 4.5	Near Real-Time	QA.3	7	3
3	AF31/AF21, 4.7/4.8	Non Real-Time	QA.4	3, Transaction Data, interactive	2
0	BE, 4.9	Best Effort	QA.5	4, Low Loss Only	1

**Table 6: Example Usage of EXP↔PHB Mappings at SP-B's ASBR**

SP-B EXP↔PHB Mapping Incoming from SP-B Domain		Generic PHB "Name"	SP-A EXP↔PHB Mapping Outgoing to MPLS-ICI		
EXP	RFC 4594 PHB		Queue	Y.1541 QoS Class, Guidance "PHB"	EXP
5	EF, 3.2	Real-Time	QA.1	6	5
4	AF41, 4.5	Near Real-Time	QA.2	7	3
3	AF31, 4.7	Non Real-Time 1	QA.3	3, Transaction Data, interactive	2
2	AF21, 4.8	Non Real-Time 2	QA.3	3, Transaction Data, interactive	2
0	BE, 4.9	Best Effort	QA.4	4, Low Loss Only	1

SP-A EXP↔PHB Mapping Incoming from MPLS-ICI		Generic PHB "Name"	SP-B EXP↔PHB Mapping Outgoing to SP-B Domain		
EXP	Y.1541 QoS Class, Guidance "PHB"		Queue	RFC 4594 PHB	EXP
5	2, Signaling/ 6	Real Time	QB.1	EF, 4.1	5
3	7	Near Real-Time	QB.2	AF41, 4.5	4
2	3, Transaction Data, interactive	Non Real Time 1	QB.3	AF31, 4.7	3
		Non Real Time 2	QB.4	AF21, 4.8	2
1	4, Low Loss Only	Best Effort	QB.5	BE, 4.9	0

In this example, the service providers are more closely aligned in terms of the generic PHB names and the EXP-bit assignments than in the previous example. However, there are some differences and the EXP↔Mapping functions are still needed. Both SP-A and SP-B have five PHBs. There is a many to one mapping in SP-A's ASBR where packets in ITU QoS classes 2 and 6 destined for SP-B are mapped to IETF PHB EF, and a null mapping for network control since SP-B has no such generic PHB. There is a many to one mapping in SP-B's ASBR where AF31 and AF21 are both mapped to ITU QoS class 3 in SP-A, and a null mapping since SP-A has support for only one non real time generic PHB while SP-B supports two non real time generic PHBs.

## II.4 Example of Canonical Forwarding Behavior Mapping

Appendix VI of [Y.1541] provides the following mapping between the two IP transfer capabilities defined in [Y.1221] and IETF Diffserv PHB specifications [RFC2475]:

**Table VI.1/Y.1541 – Association of Y.1541 QoS classes with Y.1221 transfer capabilities and differentiated services PHBs**

Y.1221 transfer capabilities	Associated DiffServ PHBs	IP QoS class	Remarks
Best-effort (BE)	Default	Unspecified QoS class 5	A legacy IP service, when operated on a lightly loaded network may achieve a good level of IP QoS.
Delay-sensitive Statistical Bandwidth (DSBW)	AF	QoS classes 2, 3, 4	The IPLR objective only applies to the IP packets in the higher priority levels of each AF class. The IPTD applies to all packets.
Dedicated Bandwidth (DBW)	EF	QoS classes 0 and 1	

Service providers connected via an MPLS-ICI could adopt this common, canonical definition of a one-one mapping of three forwarding behaviors. However, they could still use different EXP bits to represent them and a remarking would be required on the MPLS-ICI if they did so. A mapping of EXP bits to a local forwarding behavior is still also required in this case.

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